

On the Limits of Sampling-Based Reachability: Geometry, Dynamics, and Sample Complexity

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1 **Abstract:** Reachability analysis is central to safety-critical control, robotics,
2 and neural network verification, but classical computational methods, such as
3 Hamilton–Jacobi reachability and set propagation, scale poorly with state dimen-
4 sion. Sampling-based methods have emerged as a promising alternative, often pro-
5 viding finite-sample guarantees that bound the probability-mass left uncovered.
6 However, an explicit account of how the geometry of the initial set, the dynamics,
7 and the sampling law affect the accuracy of the estimator is not fully available
8 in the literature. We study this by casting sampling-based reachable-set recovery
9 as geometric support estimation over a family of problems specified by an initial
10 set, its dynamics, and a sampling law. First, we identify two regularity proper-
11 ties, positive reach of the initial set’s complement and Lipschitz continuity of the
12 dynamics, that together make recovery well-posed: a probability-mass coverage
13 guarantee can be upgraded to accuracy r in Hausdorff distance. Second, we bound
14 the resulting sample complexity: recovery is achievable with $\tilde{O}((e^{3LT}/r)^n)$ sam-
15 ples, exponential in both the state dimension and the time horizon. Third, we
16 show that neither can be removed: an almost-matching minimax lower bound
17 of $\Omega((e^{LT}/r)^n)$ holds for every estimator, so the exponential dependence on di-
18 mension and the degradation over the horizon are both intrinsic, not artifacts of
19 a particular method. Experiments on nonlinear systems confirm that adversarial
20 sampling improves constants but not the scaling.

21 **Keywords:** Reachability analysis, sample complexity, positive reach

22 1 Introduction

23 Reachability analysis is a fundamental tool for certifying the safety of dynamical and learning-
24 enabled systems [1, 2, 3, 4]. Classical computational methods, based on Hamilton–Jacobi reach-
25 ability [5, 6] or set propagation [7, 8], give strong deterministic guarantees but scale poorly with
26 state dimension. This barrier is especially pronounced in modern learning-enabled systems, where
27 neural-network controllers and learned dynamics put classical methods out of reach. To improve
28 scalability, sampling-based methods approximate reachable sets by propagating finitely many initial
29 states through the dynamics [9, 10, 11, 12]. They are model-agnostic, parallelizable, and easy to
30 implement, applicable to general nonlinear and learning-based systems. These methods are part of a
31 broader rise of sampling-based control and optimization for high-dimensional systems [13, 14, 15],
32 made practical by modern parallel hardware.

33 The accompanying finite-sample guarantees, however, are limited. They either bound the
34 probability-mass the estimate fails to cover [16, 17, 18], permitting an estimate to miss a thin, spa-
35 tially separated region without violating the bound, or they measure error in Hausdorff distance
36 against the convex hull of the reachable set [19, 20], which over-approximates nonconvex sets by

All codes and relevant materials are available at: <https://github.com/AnonyDreamer/reachapprox>

37 filling holes and merging disconnected components. What remains missing is a characterization
 38 of Hausdorff recovery of the actual set, and an account of how the sample cost depends on the
 39 geometry of the initial set, the dynamics, and the sampling law. Posed this way, the question be-
 40 comes one of geometric support estimation under flow distortion, a perspective with a rich classical
 41 theory [21, 22, 23] that has not yet been brought to bear on the reachability setting.

42 **Contributions.** To fill this gap, we study these questions for autonomous systems $\dot{x} = F(x)$,
 43 recovering the reachable set $S_T = \mathcal{R}_T(S_0)$ from endpoint samples of trajectories initialized in S_0 .
 44 Casting recovery as geometric support estimation over a family of problems specified by the initial
 45 set, its dynamics, and a sampling law, we make the following contributions:

- 46 1. **Well-posedness.** We identify two regularity properties, positive reach of the initial set’s comple-
 47 ment (r_0) and Lipschitz continuity of the dynamics (L), under which a probability-mass coverage
 48 guarantee can be upgraded to Hausdorff-accurate recovery of the reachable set.
- 49 2. **Sample complexity.** We give a random-covering upper bound showing that $\tilde{\mathcal{O}}((e^{3LT}/r)^n)$
 50 endpoint samples suffice for r -accurate inner coverage of the reachable set. Together with an
 51 estimator-dependent outer-deviation condition, this yields a Hausdorff guarantee.
- 52 3. **Unavoidability.** We prove an almost-matching minimax lower bound: any estimator, under any
 53 sampling law, requires $\Omega((e^{LT}/r)^n)$ samples on some instance, so the dimension dependence
 54 and the horizon degradation are intrinsic, not artifacts of a method.
- 55 4. **Numerical experiments.** On a 2D non-Lipschitz example outside our assumptions and a multi-
 56 link robot arm in MuJoCo, the empirical sample complexity tracks the predicted dimension-
 57 dependent scaling, and adversarial sampling improves the constants but not the scaling.

58 **Closely related work.** Our analysis is closely related to that of Lew et al. [20], who derive a finite-
 59 sample Hausdorff guarantee for a sampling-based estimator of the convex hull of the reachable set
 60 under essentially the same problem family as ours: the complement of the initial set is r -convex,
 61 the vector field is Lipschitz, and the sampling law admits a Lebesgue density lower bound. Our
 62 positive-reach assumption $\text{reach}(S_0^c) \geq r_0$ (Definition 1) is strictly stronger and implies their r -
 63 convexity. We go further in three directions: (i) we study Hausdorff recovery of the actual endpoint
 64 support, rather than the convex hull of the reachable set; (ii) we make the horizon dependence
 65 explicit, rather than absorbing it into a single Lipschitz constant of the flow; and (iii) we certify
 66 the rate with a matching minimax lower bound that holds for every sampling law. The r^{-n} rate is
 67 classical in geometric support estimation [21, 22, 23] and is known to be minimax-optimal under
 68 positive-reach conditions [24]; in our setting it persists under flow distortion with an explicit e^{nLT}
 69 amplification. A more extensive discussion of adjacent literatures, including scenario optimization,
 70 conformal prediction, and Lipschitz-based bloating, is provided in Appendix A.

71 **Notation.** For a Lebesgue measurable set $A \subseteq \mathbb{R}^n$, we write $A^c := \mathbb{R}^n \setminus A$ for its complement,
 72 \bar{A} for its closure, and ∂A for its boundary. Let $|A| := \lambda_n(A)$, where λ_n is the n -dimensional
 73 Lebesgue measure. For $x \in \mathbb{R}^n$, $r > 0$, and nonempty sets $A, B \subseteq \mathbb{R}^n$, define $\mathcal{B}_r(x) := \{y \in$
 74 $\mathbb{R}^n : \|y - x\| \leq r\}$, $d(x, A) := \inf_{a \in A} \|x - a\|$, and $\mathcal{B}_r(A) := \{x \in \mathbb{R}^n : d(x, A) \leq r\}$. The
 75 Hausdorff distance between A and B is $d_H(A, B) := \max\{\sup_{a \in A} d(a, B), \sup_{b \in B} d(b, A)\}$. Let
 76 $\omega_n := |\mathcal{B}_1(0)|$ denote the volume of the unit ball in \mathbb{R}^n . We use $A \oplus B := \{a + b : a \in A, b \in B\}$,
 77 $A \ominus B := \{x \in \mathbb{R}^n : x + B \subseteq A\}$, and $A \circ B := (A \ominus B) \oplus B$ for Minkowski sum, erosion, and
 78 opening, respectively. For a probability measure P , $\text{supp}(P)$ denotes its support, and for probability
 79 measures P, Q , $D_{\text{KL}}(P||Q)$ denotes the Kullback–Leibler divergence.

80 2 Problem Formulation

81 We consider an autonomous continuous-time dynamical system $\dot{x} = F(x)$, where $x \in \mathcal{X} \subseteq \mathbb{R}^n$ and
 82 $F : \mathcal{X} \rightarrow \mathbb{R}^n$. Given an initial condition $x_0 \in \mathcal{X}$, let $\varphi(t, x_0)$ denote the solution at time $t \geq 0$. For
 83 an initial set $S_0 \subseteq \mathcal{X}$ and a time horizon $T > 0$, the forward reachable set is

$$\mathcal{R}_T(S_0) := \{\varphi(T, x_0) \in \mathcal{X} : x_0 \in S_0\}. \quad (1)$$

84 When S_0 and T are known, we write $S_T := \mathcal{R}_T(S_0)$. This paper studies the problem of approximat-
 85 ing S_T from finitely many endpoint samples. Specifically, suppose that initial states X_1, \dots, X_N are
 86 sampled independently from a probability distribution P_0 supported on $\overline{S_0}$, and the corresponding
 87 endpoints are $Y_i := \varphi(T, X_i), i = 1, \dots, N$. These endpoints are i.i.d. samples from the push-
 88 forward distribution $P_T := (\varphi(T, \cdot))_{\#} P_0$, whose support is $\overline{S_T}$. A sampling-based reachable-set
 89 estimator is a measurable set-valued map $\widehat{S}_N = \widehat{S}_N(Y_1, \dots, Y_N) \subseteq \mathbb{R}^n$ that approximates the
 90 reachable set S_T .

91 Many existing sampling-based and learning-based approaches provide probabilistic coverage guar-
 92 antees. Informally, such results show that, with probability at least $1 - \delta$, the estimated set misses
 93 at most ε probability-mass under the endpoint distribution P_T . A typical guarantee takes the form
 94 $P_T(\overline{S_T} \setminus \widehat{S}_N) \leq \varepsilon$ with probability larger than $1 - \delta$ over the sampled endpoints and the joint dis-
 95 tribution P_T^N of (Y_1, \dots, Y_N) [16, 18]. However, probability-mass accuracy alone does not imply
 96 geometric accuracy. An estimator may miss only a small amount of probability-mass while still
 97 incurring a large Hausdorff error. This can happen when the missed region has small probability but
 98 is geometrically far from the estimated set, or when the reachable set contains thin cusps, spikes, or
 99 low-density regions. In safety and verification problems, such geometric errors are important: miss-
 100 ing a small-probability but spatially significant region may still lead to an incorrect reachable-set
 101 certificate.

102 Motivated by this gap, our goal is to understand how many endpoint samples are necessary and
 103 sufficient for \widehat{S}_N to approximate S_T in Hausdorff distance. To be more specific, for a target accuracy
 104 $r > 0$ and confidence level $1 - \delta$, we seek conditions and sample size N under which:

$$P_T^N \left(d_H(S_T, \widehat{S}_N) \leq r \right) \geq 1 - \delta. \quad (2)$$

105 This objective is stronger than controlling the probability-mass of the missed region. It requires
 106 every point of the true reachable set to lie within distance r of the estimator, and conversely requires
 107 the estimator not to extend too far away from the true reachable set. This property is arguably
 108 necessary for safety-critical circumstances.

109 The central question is therefore: when can a probability-mass guarantee be converted into a
 110 Hausdorff-distance guarantee? Such a conversion is impossible without additional structure. In-
 111 deed, if there exists a point $x \in S_T$ such that $\mathcal{B}_r(x) \cap \overline{S_T}$ has arbitrarily small P_T -mass, then an
 112 estimator may miss this entire r -scale neighborhood while still satisfying a small probability error
 113 bound. In that case, the probability error is small, but the Hausdorff distance from S_T to \widehat{S}_N is at
 114 least of order r . Thus, to make (ε, δ) -type probability guarantees geometrically meaningful, one
 115 needs a lower bound on the probability-mass of local geometric neighborhoods.

116 Figure 1 illustrates this phen-
 117 omenon on the autonomous
 118 system $\dot{y} = 0, \dot{x} = x^2$. Start-
 119 ing from two initial sets with
 comparable geometric scale,
 a disk and a nonconvex
 star, we uniformly sample
 initial conditions, propagate
 the samples, and construct
 support estimators from the
 endpoints. The flow map
 expands regions with large
 x , causing geometric error
 to increase with time, i.e.
 the Hausdorff distance grows
 super-exponentially as the
 time horizon increases, reflect-

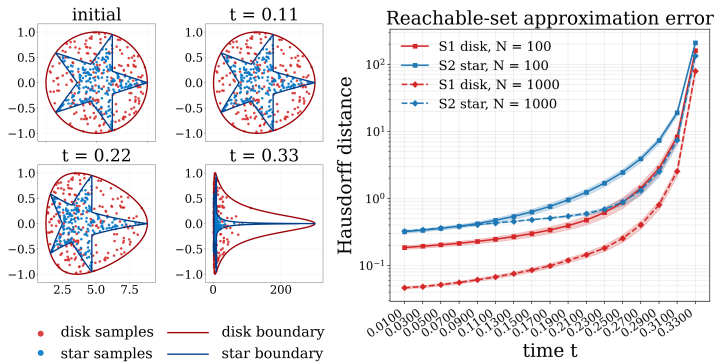


Figure 1: Reachable-set approximation under the autonomous dynamics $\dot{y} = 0, \dot{x} = x^2$.

117 Moreover, for the same
 118 sampling budget and comparable geometric scale, the star-shaped set produces larger Hausdorff
 119 error than the disk. This behavior is not explained by probability-mass alone. Rather, it reflects

120 two separate effects. First, the dynamics can distort local neighborhoods and probability-mass over
121 time. Second, the geometry of the initial set determines whether every small spatial neighborhood
122 contains enough volume, and hence enough probability-mass, to be detected from samples. This
123 discussion motivates the following problem class. We require the flow to be well posed over the finite
124 horizon of interest and Lipschitz with respect to initial conditions, so that local neighborhoods can
125 be related through the pushforward dynamics. We also impose a geometric regularity condition on
126 the initial set to exclude thin degeneracies. Throughout, for each finite horizon T considered below,
127 we assume that $\varphi(t, \cdot)$ is well defined and unique on the relevant initial states for all $t \in [0, T]$; this
128 is a standing well-posedness convention, separate from the statistical problem class. To state the
129 geometric condition, we first recall the notion of reach.

130 **Definition 1** (Reach [25]). *Let $S \subset \mathbb{R}^n$ be nonempty and closed. For $x \in \mathbb{R}^n$, define $d(x, S) :=$
131 $\inf_{y \in S} \|x - y\|$ and $\Pi_S(x) := \{y \in S : \|x - y\| = d(x, S)\}$. The reach of S is defined as follows:*

$$\text{reach}(S) := \sup \{r \geq 0 : \Pi_S(x) \text{ is a singleton whenever } d(x, S) < r\}.$$

132 We now collect the dynamical and geometric requirements into a single problem family.

133 **Definition 2** (Problem family). *Fix constants $L > 0$ and $r_0 > 0$. Let \mathcal{F}_{L, r_0} denote the class of
134 triples (S_0, F, P_0) such that:*

- 135 1. $S_0 \subset \mathcal{X}$ is a bounded open set with $0 < |S_0| < \infty$;
- 136 2. $\text{reach}(S_0^c) \geq r_0$, where $S_0^c := \mathbb{R}^n \setminus S_0$;
- 137 3. $F : \mathcal{X} \rightarrow \mathbb{R}^n$ is L -Lipschitz on \mathcal{X} , i.e., $\|F(x) - F(y)\| \leq L\|x - y\|, \forall x, y \in \mathcal{X}$;
- 138 4. the initial distribution P_0 satisfies $\text{supp}(P_0) = \overline{S_0}$.

139 The problem family allows arbitrary finite positive volume. Throughout the paper, we use the nor-
140 malization $|S_0| = 1$ only to fix the spatial scale and simplify constants. This normalization is
141 not essential and can be replaced by any other fixed finite volume, with the corresponding rescal-
142 ing of the reach parameter. This family gives a local thickness property: by [26, Lemma 4.8],
143 $\overline{S_0} = \overline{S_0} \circ \mathcal{B}_{r_0}(0)$, positive reach of the complement rules out inward cusps and thin degeneracies,
144 yielding lower bounds of the form $|\mathcal{B}_\rho(x) \cap \overline{S_0}| \gtrsim \rho^n$ for some small $\rho > 0$.¹ This local volume
145 bound is the bridge from probability to geometry. After accounting for Lipschitz flow distortion and
146 endpoint density lower bounds, it ensures that missing an r -scale region of the reachable set incurs
147 non-negligible probability mass, which enables Hausdorff accuracy guarantees.

148 3 Reachable Set Approximation under Regular Geometry

149 Under the regularity conditions introduced above, we study Hausdorff approximation of the reach-
150 able set from finite endpoint samples. The key ingredient is a local endpoint probability lower
151 bound, obtained by combining positive reach, an initial density lower bound, and Lipschitz flow
152 distortion. This bound yields an inner-coverage sample-complexity result for any estimator contain-
153 ing all endpoint samples; a full Hausdorff guarantee follows after adding an estimator-dependent
154 outer-deviation condition. We then prove that the r^{-n} dependence is information-theoretically un-
155 avoidable.

156 3.1 Local Mass Bounds under Flow Distortion

157 We first establish local mass lower bounds for the reachable set. Positive reach gives a static vol-
158 ume lower bound near the initial support, and the probability density lower bound converts it into
159 sampling probability.

160 **Proposition 1** (Lower Lebesgue mass bound for the initial set). *Let $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$. Then, for
161 every $x \in \overline{S_0}$ and every $r \in (0, r_0]$, we have $|\mathcal{B}_r(x) \cap \overline{S_0}| \geq \omega_n 2^{-n} r^n$.*

¹ $A \gtrsim B$ means $A \geq cB$ for a constant $c > 0$ independent of the scale of B .

162 *Proof.* See Appendix B.1. □

163 **Assumption 1** (Initial density lower bound). *The sampling distribution P_0 is absolutely continuous*
 164 *with respect to Lebesgue measure on $\overline{S_0}$, with density p_0 . Moreover, there exists $\underline{p}_0 > 0$ such that*
 165 *$p_0(x) \geq \underline{p}_0$, for a.e. $x \in \overline{S_0}$.*

166 Combining Proposition 1 with Assumption 1 gives an initial local probability lower bound, making
 167 explicit the roles of r_0 and \underline{p}_0 .

168 **Corollary 1** (Initial local probability lower bound). *Suppose $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$ and Assumption 1*
 169 *holds. Then, for every $x \in \overline{S_0}$ and every $r \in (0, r_0]$,*

$$P_0(\mathcal{B}_r(x) \cap \overline{S_0}) \geq \underline{p}_0 \omega_n 2^{-n} r^n.$$

170 To apply this local-mass argument to the reachable set at time T , we do not track the reach parameter
 171 itself. Instead, we propagate the local thickness implied by $\text{reach}(S_0^c) \geq r_0$ through the Lipschitz
 172 flow and combine it with the propagated density lower bound.

173 **Proposition 2** (Propagation of local Lebesgue mass). *Let $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$. Then, for every*
 174 *$x \in \overline{S_T}$ and every $r \in (0, e^{LT} r_0]$, $|\mathcal{B}_r(x) \cap \overline{S_T}| \geq \omega_n 2^{-n} e^{-2nLT} r^n$.*

175 *Proof.* See Appendix B.3. □

176 **Proposition 3** (Propagation of density lower bounds). *Suppose $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$ and Assump-*
 177 *tion 1 holds. Then, P_T is absolutely continuous with respect to Lebesgue measure on $\overline{S_T}$. Moreover,*
 178 *if p_T denotes its density, then $p_T(y) \geq \underline{p}_0 e^{-nLT}$, for a.e. $y \in \overline{S_T}$. In particular, at the terminal time*
 179 *T , $\underline{p}_T := \underline{p}_0 e^{-nLT}$ is the sample density lower bound for the reachable set at time T .*

180 *Proof.* See Appendix B.3. □

181 Combining Proposition 2 with Proposition 3, we obtain the local endpoint probability lower bound
 182 used in the sampling-covering argument in the next subsection.

183 **Corollary 2** (Endpoint local probability lower bound). *Suppose $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$ and As-*
 184 *sumption 1 holds. Define $\Gamma_T(r) := \underline{p}_0 \omega_n 2^{-n} e^{-3nLT} r^n$. Then, for every $x \in \overline{S_T}$ and every*
 185 *$r \in (0, e^{LT} r_0]$,*

$$P_T(\mathcal{B}_r(x) \cap \overline{S_T}) \geq \Gamma_T(r).$$

186 *Proof.* See Appendix B.3. □

187 3.2 Sample Complexity Bounds

188 Corollary 2 shows that every r -neighborhood centered on the reachable set carries at least $\Gamma_T(r)$
 189 endpoint probability mass. We now use this local probability lower bound to quantify how many
 190 endpoint samples are sufficient to cover the reachable set in the inner direction. The argument is
 191 estimator-agnostic: it only requires that the estimator contains all sampled endpoints.

192 **Theorem 1** (Sampling upper bound for inner coverage). *Let $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$, and Y_1, \dots, Y_N*
 193 *be i.i.d. samples from P_T . Suppose that Assumption 1 holds and the estimator \widehat{S}_N contains all*
 194 *endpoint samples, i.e., $Y_i \in \widehat{S}_N$ for all $i = 1, \dots, N$. Fix $0 < r \leq 2e^{LT} r_0$. If*

$$N \geq \frac{2^{2n} e^{3nLT}}{\underline{p}_0 \omega_n r^n} \left[\log \left(\frac{2^{3n} e^{nLT}}{\omega_n r^n} \right) + \log \frac{1}{\delta} \right],$$

195 *then, with probability at least $1 - \delta$, $\sup_{x \in \overline{S_T}} d(x, \widehat{S}_N) \leq r$.*

196 *Proof.* See Appendix B.4. □

197 **Corollary 3** (Hausdorff bound with outer deviation). *Under the conditions of Theorem 1, suppose*
 198 *in addition that the estimator satisfies the outer-deviation condition $\widehat{S}_N \subseteq \mathcal{B}_\eta(S_T)$ for some $\eta > 0$.*
 199 *Then, with probability at least $1 - \delta$, $d_H(S_T, \widehat{S}_N) \leq \max\{r, \eta\}$.*

200 *Proof.* See Appendix B.4. □

201 **Remark 1** (Estimator-dependent outer deviation). *The parameter η measures the outer deviation*
 202 *of the estimator where $\widehat{S}_N \subseteq \mathcal{B}_\eta(S_T)$ is equivalent to $\sup_{z \in \widehat{S}_N} d(z, S_T) \leq \eta$. In Corollary 3, the*
 203 *missed-mass condition controls the directed error from S_T to \widehat{S}_N , while η controls the opposite*
 204 *direction. Its value is estimator-dependent: $\eta = 0$ for inner estimators, $\eta = h$ for a union-of-balls*
 205 *estimator with radius h , and for a convex-hull estimator it is governed by the convexification error*
 206 *$\sup_{z \in \text{conv}(S_T)} d(z, S_T)$. For a Christoffel-function estimator $\widehat{S}_N^{\text{Chr}} := \{z \in \mathcal{X} : \kappa_N(z) \leq \tau_N\}$, the*
 207 *corresponding value is $\eta_{\text{Chr}, N} := \sup_{\kappa_N(z) \leq \tau_N} d(z, S_T)$.*

208 We next show that the r^{-n} dependence is not a proof artifact, but an intrinsic limitation of sampling-
 209 based reachable-set approximation.

210 **Theorem 2** (Minimax Hausdorff lower bound). *Let $C := 3^n \omega_n$. For any accuracy $0 < r \leq$*
 211 *$2^{-\frac{n+1}{n}} e^{LT} r_0$ and any confidence level $1 - \delta$ with $\delta \in (0, 1)$, if $N < \frac{e^{nLT}}{2^{n+1} C r^n} \log \frac{1}{4\delta}$ and $r_0 \leq C^{-\frac{1}{n}}$,*
 212 *then for every estimator \widehat{S}_N , there exists an instance $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$ such that*

$$P_T^N(d_H(S_T, \widehat{S}_N) > r) > \delta.$$

213 *Proof.* See Appendix B.4. □

214 Together, the bounds show that the intrinsic dependence on the accuracy parameter is r^{-n} . The up-
 215 per bound gives inner-coverage achievability up to logarithmic and flow-dependent factors, while the
 216 lower bound shows that $\Omega(e^{nLT} r^{-n} \log(1/\delta))$ samples are unavoidable in the worst case. Table 1
 217 summarizes the comparison.

Table 1: Comparison between the upper and lower bounds.

	Upper bound	Lower bound
Type	Achievability	Minimax Hausdorff impossibility
Rate	$\mathcal{O}(e^{3nLT} r^{-n} [\log(1/\delta) + \log(e^{nLT} r^{-n})])$	$\Omega(e^{nLT} r^{-n} \log(1/\delta))$

218 4 Algorithm Design

219 The r^{-n} dependence in the lower bound highlights the intrinsic coverage burden of Hausdorff
 220 reachable-set approximation. This motivates the practical question studied in our experiments: un-
 221 der a fixed sample budget, can sampling be allocated more effectively than uniform sampling so
 222 as to improve geometric coverage and reduce the empirical burden of this curse of dimensionality?
 223 Accordingly, we focus on the sampling stage rather than on the downstream set representation. In
 224 the autonomous setting, samples are selected only through their initial conditions in S_0 . For each
 225 sampled initial condition $X_j \in S_0$, we write $Y_j := \varphi(T, X_j)$ for the corresponding endpoint. Given
 226 endpoint samples $\mathcal{Y}_N = \{Y_j\}_{j=1}^N$, the reachable-set estimate is constructed as $\widehat{S}_N = C(\mathcal{Y}_N)$, where
 227 $C(\cdot)$ denotes a generic reachable-set estimator.

228 In the experiments, we compare uniform sampling with the adversarial sampling heuristic described
 229 by Lew and Pavone [19]. The purpose of using these two different sampling algorithms is to
 230 investigate whether adversarial sampling can improve finite-sample coverage by directing samples
 231 toward geometrically informative regions. We emphasize that the adversarial sampling procedure
 232 can improve finite sample coverage but cannot violate the lower bound. The worst-case Hausdorff
 233 sample complexity remains governed by the intrinsic r^{-n} coverage requirement. Implementation
 234 details of the adversarial sampling procedure (Algorithm 1) are provided in Appendix C.2. When
 235 $n_{\text{adv}} = 0$ in Algorithm 1, no adversarial update is performed, and the method reduces to the uniform
 236 sampling method used in the experiments.

237 **5 Experiments**

238 In this section, we compare uniform and adversarial sampling on two 2D systems, one non-Lipschitz
 239 and one Lipschitz, and on a closed-loop robot-arm control task with different dimensions.

240 **5.1 Adversarial Reachable Set Approximation under Non-Lipschitz Dynamics**

241 Using the system in Figure 1, we first evaluate reachable-set approximation for the non-Lipschitz
 242 autonomous dynamics. This vector field is not globally Lipschitz and exhibits finite-time blow-up,
 243 so sampling errors can be strongly amplified in the positive x -direction. We consider three initial
 244 sets centered at $(2, 0)$: a disk, an equilateral triangle, and a morphologically opened triangle, all
 245 with the same area but different boundary geometry. For each set, we approximate the reachable
 246 set by the convex hull of propagated samples and report its Hausdorff distance to a high-resolution
 247 reference reachable set over time. Solid and dashed curves denote uniform and adversarial sampling,
 248 with shaded regions showing 95% confidence intervals over 50 random seeds.

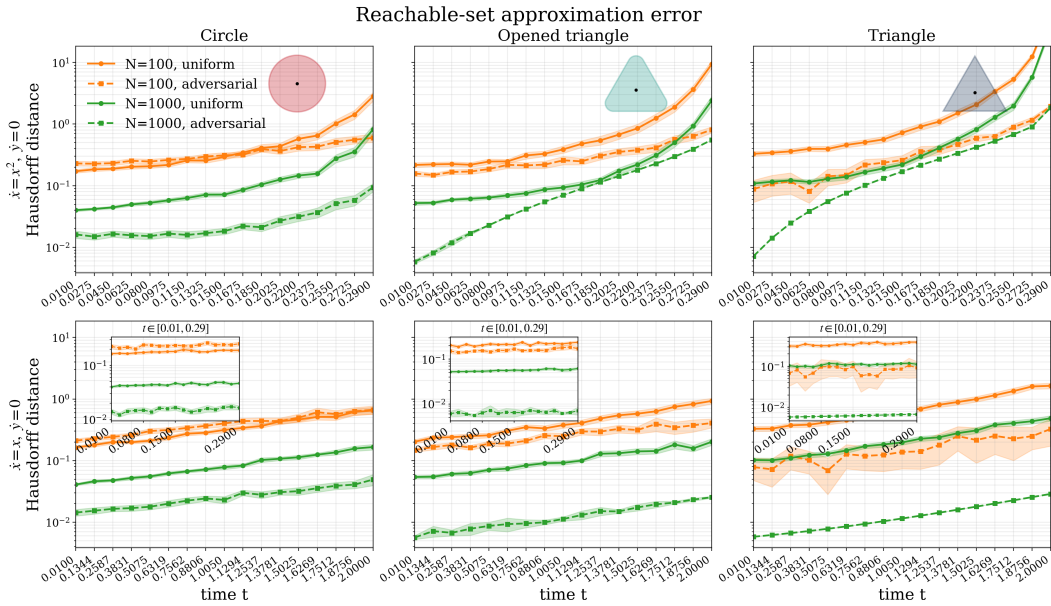


Figure 2: Hausdorff error versus time for $\dot{y} = 0, \dot{x} = x^2$ and $\dot{y} = 0, \dot{x} = x$ using different sampling methods. Solid and dashed curves denote uniform and adversarial sampling.

249 To isolate the effect of flow regularity, we also compare with the Lipschitz linear system $\dot{y} = 0, \dot{x} =$
 250 x under the same initial-set configurations. Figure 2 shows that the linear system follows the ex-
 251 pected exponential flow expansion, while the nonlinear system grows much faster due to finite-time
 252 blow-up. Across both systems, adversarial sampling reduces finite-sample error but mainly im-
 253 proves constants; the triangle and opened triangle remain harder than the disk, reflecting the effect
 254 of initial-set geometry.

255 **5.2 Intrinsic Curse of Dimensionality Cannot be Circumvented**

256 We also consider a vertical n -link robot arm tracking problem with $n \in \{2, 3, 4\}$. The state is
 257 $x = [q^\top, v^\top]^\top \in \mathbb{R}^{2n}$, where $q \in \mathbb{R}^n$ denotes joint angles and $v = \dot{q} \in \mathbb{R}^n$ denotes joint velocities.
 258 The trajectories are simulated in MuJoCo under the rigid-body dynamics $M(q)\dot{v} + C(q, v)v +$
 259 $g(q) = \tau$. Given a smooth reference trajectory $q_d(t)$, we use a non-adaptive tracking controller
 260 $\tau = M(q)\ddot{q}_d(t) + C(q, v)v + g(q) - K_p(q - q_d(t)) - K_d(v - \dot{q}_d(t))$ for control.

261 Although the nominal task is trajectory tracking, the initial state is uncertain, so safety verifica-
 262 tion requires characterizing how this uncertainty propagates through the closed-loop dynamics. We
 263 model the initial uncertainty by $S_0 = ([-\rho_q, \rho_q]^n \times [-\rho_v, \rho_v]^n) \circ \mathcal{B}_r(0), \rho_q = \rho_v = 0.1, r = 0.01$.

264 We propagate sampled initial states to $T = 1.0$ second and measure approximation error by the
 265 point-cloud directed Hausdorff distance from a dense reference terminal cloud to the convex hull of
 266 the sampled terminal states.

267 As shown in Figure 3 and Table 2, the Hausdorff error decreases with the sample size under both uni-
 form and adversarial sampling. However, the decay becomes slower as the state dimension increases: the fitted log-log slopes become less negative from dimension 4 (2-link) to dimension 8 (4-link). Adversarial sampling gives steeper slopes than uniform sampling in all dimensions, indicating better finite-sample efficiency. Nevertheless, its slopes also flatten as the dimension increases. Thus, targeted sampling improves the constants but does not remove the intrinsic dimension dependent

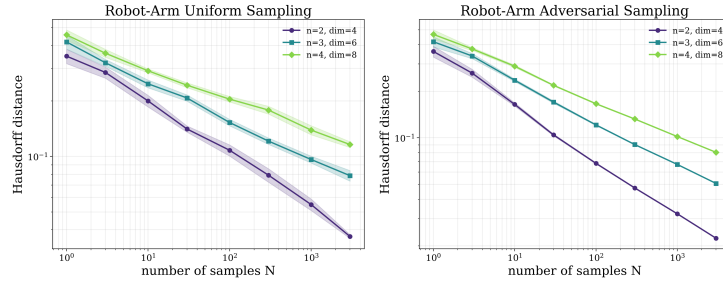


Figure 3: Hausdorff error versus sample size for robot-arm uncertainty propagation corresponding to different state dimensions under uniform sampling and adversarial sampling.

Dimension	4	6	8
Uniform sampling	-0.2806	-0.2094	-0.1662
Adversarial sampling	-0.3535	-0.2701	-0.2224

268 degradation predicted by the r^{-n} -type sample-complexity scaling. Additional implementation details and more detailed comparisons and explanation, including the sample-budget-dependent improvement of adversarial sampling, are provided in Appendix C.3.

271 6 Conclusion and Future Work

272 This paper studied the finite-sample limits of sampling-based reachable-set approximation in Haus-
 273 dorff distance. By viewing endpoint samples as observations from a pushforward distribution, we
 274 formulated reachable-set recovery as a geometric support estimation problem. We showed that,
 275 under positive-reach regularity of the initial domain and Lipschitz continuity of the autonomous
 276 flow, probability-mass coverage can be upgraded to Hausdorff accuracy. Our results give matching
 277 upper- and lower-bound perspectives. The upper bound shows that Hausdorff accuracy is attain-
 278 able once every relevant r -scale region of the reachable set receives samples. The minimax lower
 279 bound shows that the dominant r^{-n} dependence, together with the flow-induced factor e^{nLT} , is
 280 intrinsic in the worst case. Experiments further support this picture: adversarial sampling improves
 281 finite-sample performance by emphasizing geometrically informative regions, but cannot remove
 282 the intrinsic dimension-dependent scaling. A natural direction for future work is to go beyond geo-
 283 metric regularity and global Lipschitz continuity by exploiting physics-informed inductive biases in
 284 the dynamics [27], such as energy conservation [28], contraction [29], or Hamiltonian structure [30],
 285 which may reduce the effective sample complexity for physically structured systems [31].

286 7 Limitations

287 The analysis is restricted to autonomous dynamics. This includes uncontrolled systems and closed-
 288 loop systems, but not open-loop reachable sets, which are unions over a control class rather than
 289 images of S_0 under a single flow map. As a result, the injectivity, volume-distortion, and endpoint-
 290 density arguments used here do not directly apply. Moreover, endpoint mass depends on how control
 291 signals are sampled, so geometrically important regions may receive very small probability mass.

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457 **Appendix**

458 The Appendix is organized as follows. Section A provides additional related work on sampling-
459 based reachable set approximation methods. Section B presents the proofs of the theoretical results
460 in the main paper, including the lower Lebesgue-mass bound, the mass-to-Hausdorff conversion
461 argument, and sample complexity’s lower bound and upper bound. Section C reports additional
462 numerical experiments and implementation details.

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486 A Further Related Work

487 Beyond the Hausdorff-oriented works discussed in the main paper, much of the sampling-based
488 reachability literature studies weaker probabilistic notions of approximation, such as volume cover-
489 age, and violation probability, which are often characterized by scenario optimization and conformal
490 prediction. We briefly review these adjacent frameworks below. While some of them provide im-
491 portant finite-sample and high-confidence guarantees, they do not by themselves ensure geometric
492 recovery of the full reachable set. This makes them complementary to our goal of understanding
493 when probability-level guarantees can be upgraded to Hausdorff accuracy.

494 A.1 Other Sampling-based Reachable Set Approximation Methods

495 **Scenario optimization and sampling-based certificates.** Scenario optimization provides a prin-
496 ciple framework for deriving probabilistic guarantees from finitely many sampled constraints. In
497 the classical scenario approach, an infinite family of constraints is replaced by a finite set of sampled
498 constraints, leading to high-confidence bounds on the probability of constraint violation [32, 33].
499 This viewpoint has been adopted in reachability and safety verification, where sampled reachable
500 states are used to construct data-driven outer approximations or safety certificates. For instance,
501 scenario-based methods have been used to estimate reachable sets within tractable templates such
502 as boxes, ellipsoids, and other parameterized set families [34, 35]. Related work has also studied
503 reach-avoid problems for discrete-time linear time-invariant systems with additive probabilistic un-
504 certainty, using Voronoi partition-based scenario reduction to accelerate sampling-based stochastic
505 reachability computation [36]. Scenario-based techniques have also been extended to nonconvex set
506 representations [37] and model-predictive control settings [38]. These methods are attractive since
507 they reduce reachable-set estimation or safety certification to finite-dimensional optimization prob-
508 lems with explicit confidence bounds. However, their guarantees are typically expressed in terms of
509 violation probability or marginal coverage, rather than Hausdorff distance. This distinction is im-
510 portant: a set of small probability-mass can still be geometrically far from the estimator, especially
511 in the presence of low-density regions, thin components, or separated reachable branches. Thus,
512 such guarantees do not directly imply Hausdorff accuracy without additional geometric regularity
513 and local lower-mass assumptions.

514 **Set growth and Lipschitz constant-based bloating.** Another line of sampling-based reachability
515 methods estimates how sets grow under the dynamics by bounding the sensitivity of trajectories to
516 initial conditions, disturbances, or model parameters. These methods often rely on Lipschitz con-
517 stants, local discrepancy functions, finite-order Taylor approximations, or learned expansion rates to
518 inflate sampled trajectories into tubes covering nearby trajectories. For example, some algorithms
519 use smoothness of the dynamics together with Lipschitz constants to control the Lagrange remain-
520 der of Taylor approximations and thereby compute outer approximations of reachable sets [39].
521 Tools such as DryVR use discrepancy functions to quantify how the distance between trajectories
522 evolves over time, allowing finitely many simulations to certify a larger reach tube [40]. More
523 recent statistical verification methods, such as GoTube, estimate local expansion or Lipschitz-like
524 quantities to construct probabilistic tubes for continuous-depth models and neural ODEs [41, 42].
525 A key limitation of this line of work is that tight Lipschitz constants are rarely available in data-
526 driven or learning-based control systems. Computing upper bounds on the Lipschitz constants of
527 neural networks is itself an active research problem [43, 44], and such bounds can be computationally
528 expensive or overly conservative. Although scalable sampling-based procedures can estimate
529 Lipschitz-like quantities, these estimates are generally not guaranteed to be valid upper bounds.
530 These approaches are closely related to our flow-dependent analysis, since both use sensitivity of
531 the dynamics to translate finite sampling resolution into geometric coverage after propagation. Our
532 theory instead isolates the sample-complexity dependence needed for Hausdorff approximation of
533 the reachable set itself, showing how flow distortion enters through dimension-dependent factors.

534 **Conformal prediction and distribution-free coverage.** Conformal prediction offers another
535 route to finite-sample guarantees. Given a calibration dataset and a user-defined nonconformity
536 score, conformal methods compute an empirical quantile of the calibration scores to construct pre-
537 diction sets with distribution-free marginal coverage, where a future test sample is contained in the
538 conformal set with a prescribed probability [45]. Since this construction does not require a paramet-
539 ric model of the data-generating distribution, it has been widely used as an uncertainty-quantification
540 tool in learning-enabled autonomous systems, including perception, prediction, safe control, offline
541 verification, and online monitoring [46, 47, 48, 49]. For instance, in safe planning for dynamic en-
542 vironments, conformal prediction calibrates prediction regions around learned trajectory predictors
543 and incorporates them into model predictive control constraints to obtain user-specified probabilistic
544 safety guarantees [50]. In reachability and safety verification, conformal prediction is typically used
545 as a calibration layer on top of a data-driven reachable-set representation or a learned reachability
546 surrogate. It has been used to calibrate Christoffel function’s sublevel set thresholds for data-driven
547 reachability, improving sample efficiency and robustness to outliers [51]; to verify DeepReach-
548 style [52] neural HJ reachable tubes with probabilistic safety guarantees despite outlier errors [53];
549 and to inflate neural surrogate flowpipes for black-box stochastic systems using conformal residual
550 bounds [54]. Recent PCA-based variants further reduce the conservatism of this inflation and im-
551 prove scalability under distribution shift [55]. These methods are appealing because they require
552 weak distributional assumptions and provide finite-sample coverage or safety guarantees. However,
553 their guarantees remain probabilistic: a conformal set contains a future sample, trajectory, or pre-
554 diction error with prescribed probability, but need not be close to the full reachable set in Hausdorff
555 distance. It may cover most endpoint samples while missing a low-probability but spatially sep-
556 arated component. Thus, conformal prediction is complementary to our analysis, which identifies
557 sampling conditions under which probability-mass coverage can be upgraded to Hausdorff accuracy.

558 B Proofs and Auxiliary Results

559 B.1 Lower Lebesgue Mass Bound for the Initial Set

560 We prove the local thickness property used in Proposition 1. Recall and that the positive-reach
561 condition is imposed on the closed complement $S_0^c = \mathbb{R}^n \setminus S_0$.

562 **Proposition 1** (Lower Lebesgue mass bound for the initial set). Let $(S_0, F, P_0) \in \mathcal{F}_{L,r_0}$. Then, for
563 every $x \in \overline{S_0}$ and every $r \in (0, r_0]$,

$$|\mathcal{B}_r(x) \cap \overline{S_0}| \geq \omega_n 2^{-n} r^n.$$

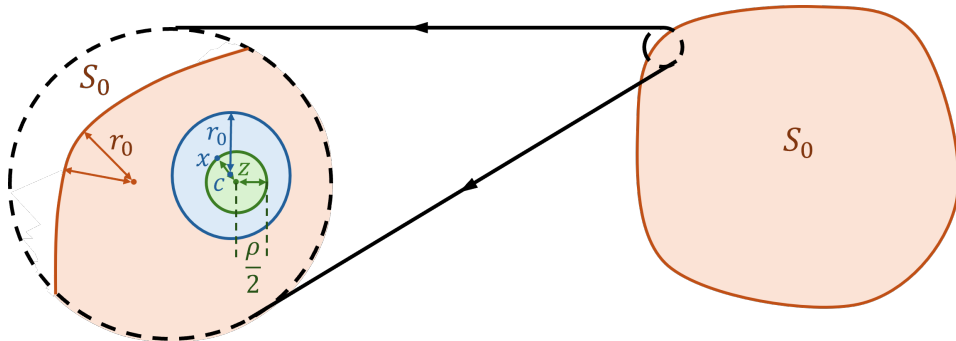


Figure 4: Local ball inclusion used in the proof of Proposition 1.

564 *Proof of Proposition 1.* Fix $x \in \overline{S_0}$ and $r \in (0, r_0]$. By the positive-reach opening property [26,
565 Lemma 4.8], the condition $\text{reach}(S_0^c) \geq r_0$ implies that $\overline{S_0}$ has an interior rolling-ball representation
566 at scale r_0 : for every $x \in \overline{S_0}$, there exists $c \in \overline{S_0}$ such that $x \in \mathcal{B}_{r_0}(c), \mathcal{B}_{r_0}(c) \subseteq \overline{S_0}$. We claim

567 that $\mathcal{B}_r(x) \cap \overline{S_0}$ contains a Euclidean ball of radius $r/2$. This immediately implies the desired lower
 568 bound. Let $d := \|x - c\|$. If $x = c$, then since $r \leq r_0$,

$$\mathcal{B}_r(x) \subseteq \mathcal{B}_{r_0}(c) \subseteq \overline{S_0}.$$

569 Therefore,

$$|\mathcal{B}_r(x) \cap \overline{S_0}| = \omega_n r^n \geq \omega_n 2^{-n} r^n.$$

570 Now suppose $x \neq c$. Define $\nu := \frac{c-x}{\|c-x\|}$, $z := x + \frac{r}{2}\nu$. Then $\|z - x\| = r/2$, and hence $\mathcal{B}_{r/2}(z) \subseteq$
 571 $\mathcal{B}_r(x)$. It remains to show that $\mathcal{B}_{r/2}(z) \subseteq \mathcal{B}_{r_0}(c)$. Take any $w \in \mathcal{B}_{r/2}(z)$. If $d \geq r/2$, then
 572 $\|z - c\| = d - r/2$, and therefore

$$\|w - c\| \leq \|w - z\| + \|z - c\| \leq \frac{r}{2} + \left(d - \frac{r}{2}\right) = d \leq r_0.$$

573 If $d < r/2$, then $\|z - c\| = r/2 - d$, and hence

$$\|w - c\| \leq \|w - z\| + \|z - c\| \leq \frac{r}{2} + \left(\frac{r}{2} - d\right) = r - d \leq r \leq r_0.$$

574 Thus, in both cases, $w \in \mathcal{B}_{r_0}(c)$. Since w was arbitrary,

$$\mathcal{B}_{r/2}(z) \subseteq \mathcal{B}_r(x) \cap \mathcal{B}_{r_0}(c) \subseteq \mathcal{B}_r(x) \cap \overline{S_0}.$$

575 Taking Lebesgue measure gives

$$|\mathcal{B}_r(x) \cap \overline{S_0}| \geq |\mathcal{B}_{r/2}(z)| = \omega_n \left(\frac{r}{2}\right)^n = \omega_n 2^{-n} r^n.$$

576 This proves the claim. □

577 B.2 Closure of the Reachable Set

578 We first record a simple topological fact used to identify the support of the endpoint distribution.

579 **Lemma 1** (Reachability commutes with closure). *Suppose $S_0 \subset \mathbb{R}^n$ is bounded and the finite-time*
 580 *flow map $\varphi(T, \cdot)$ is continuous on $\overline{S_0} = \overline{S_0}$. Then*

$$\mathcal{R}_T(\overline{S_0}) = \varphi(T, \overline{S_0}) = \overline{\varphi(T, S_0)} = \overline{\mathcal{R}_T(S_0)}.$$

581 *In particular, if $S_T := \mathcal{R}_T(S_0)$, then $\mathcal{R}_T(\overline{S_0}) = \overline{S_T}$.*

582 *Proof.* Since S_0 is dense in $\overline{S_0}$ and $\varphi(T, \cdot)$ is continuous on $\overline{S_0}$, we have $\varphi(T, \overline{S_0}) \subseteq \overline{\varphi(T, S_0)}$.
 583 Conversely, since S_0 is bounded, $\overline{S_0}$ is compact. Therefore $\varphi(T, \overline{S_0})$ is compact, hence closed.
 584 Moreover, $\varphi(T, S_0) \subseteq \varphi(T, \overline{S_0})$, so

$$\overline{\varphi(T, S_0)} \subseteq \varphi(T, \overline{S_0}).$$

585 Combining the two inclusions gives the claim. □

586 B.3 Volume Distortion and Density Propagation

587 We next prove the flow-regularity and density-propagation statements used in the main text.
 588 Throughout this subsection, $\overline{S_0} = \overline{S_0}$ and $\overline{S_T} := \mathcal{R}_T(\overline{S_0}) = \overline{\mathcal{R}_T(S_0)}$ by Lemma 1. We begin
 589 with the basic distortion properties of the finite-time flow. These estimates show that the flow is
 590 bi-Lipschitz on the relevant support and that Lebesgue volume cannot shrink by more than a factor
 591 e^{-nLT} .

592 **Proposition 4** (Lipschitz regularity and geometry distortion of the flow). *Let $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$.*
 593 *Then we have:*

594 1. *the map $x \mapsto \varphi(T, x)$ from $\overline{S_0}$ to $\overline{S_T}$ is injective;*

595 2. for all

$$x, y \in \overline{S_0}, \|\varphi(T, x) - \varphi(T, y)\| \leq e^{LT} \|x - y\|;$$

596

597 3. for all

$$x, y \in \overline{S_0}, \|\varphi(T, x) - \varphi(T, y)\| \geq e^{-LT} \|x - y\|,$$

598 equivalently, the inverse map $\varphi(T, \cdot)^{-1} : \overline{S_T} \rightarrow \overline{S_0}$ is e^{LT} -Lipschitz;

599 4. for every Lebesgue measurable set

$$A \subseteq \overline{S_0}, |\mathcal{R}_T(A)| \geq e^{-nLT} |A|. \quad (3)$$

600

601 *Proof of Proposition 4.* We prove the four claims in such an order: 2 \rightarrow 1 \rightarrow 3 \rightarrow 4. First, we
 602 prove the upper Lipschitz bound. Fix $x, y \in \overline{S_0}$. Since $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$, the vector field F is
 603 L -Lipschitz on \mathcal{X} . Hence, for all $s \in [0, T]$,

$$\|F(\varphi(s, x)) - F(\varphi(s, y))\| \leq L \|\varphi(s, x) - \varphi(s, y)\|.$$

604 Using the integral form of the flow,

$$\varphi(s, x) - \varphi(s, y) = x - y + \int_0^s [F(\varphi(\sigma, x)) - F(\varphi(\sigma, y))] d\sigma.$$

605 Taking norms gives

$$\|\varphi(s, x) - \varphi(s, y)\| \leq \|x - y\| + L \int_0^s \|\varphi(\sigma, x) - \varphi(\sigma, y)\| d\sigma.$$

606 By Gronwall's inequality [56, Corollary 3.17],

$$\|\varphi(s, x) - \varphi(s, y)\| \leq e^{Ls} \|x - y\|, \quad s \in [0, T].$$

607 Taking $s = T$ proves

$$\|\varphi(T, x) - \varphi(T, y)\| \leq e^{LT} \|x - y\|.$$

608 Next, we prove injectivity. Suppose that $\varphi(T, a) = \varphi(T, b)$ for some $a, b \in \overline{S_0}$. Define the reverse-
 609 time trajectories

$$\eta_a(s) := \varphi(T - s, a), \quad \eta_b(s) := \varphi(T - s, b), \quad s \in [0, T].$$

610 Then $\eta_a(0) = \eta_b(0)$, and both trajectories solve the reverse-time equation

$$\dot{\eta}(s) = -F(\eta(s)).$$

611 Since F is L -Lipschitz, the reverse-time vector field $-F$ is also L -Lipschitz. By uniqueness of
 612 solutions, $\eta_a(s) = \eta_b(s)$ for all $s \in [0, T]$. Evaluating at $s = T$ gives $a = b$. Thus $x \mapsto \varphi(T, x)$ is
 613 injective on $\overline{S_0}$.

614 We now prove the lower Lipschitz bound. Fix $x, y \in \overline{S_0}$ and set

$$p := \varphi(T, x), \quad q := \varphi(T, y).$$

615 Apply the upper Lipschitz estimate to the reverse-time dynamics. Since the reverse flow from p and
 616 q over time T returns to x and y , respectively, we obtain

$$\|x - y\| \leq e^{LT} \|p - q\| = e^{LT} \|\varphi(T, x) - \varphi(T, y)\|.$$

617 Equivalently,

$$\|\varphi(T, x) - \varphi(T, y)\| \geq e^{-LT} \|x - y\|.$$

618 This also shows that the inverse map $\varphi(T, \cdot)^{-1} : \overline{S_T} \rightarrow \overline{S_0}$ is e^{LT} -Lipschitz.

619 It remains to prove the volume distortion bound. By [57, Theorem 2.8], we know that the inverse
620 flow map $\varphi(T, \cdot)^{-1} : \overline{S_T} \rightarrow \overline{S_0}$, which is e^{LT} -Lipschitz, gives $|\varphi(T, \cdot)^{-1}(B)| \leq e^{nLT}|B|, \forall B \subseteq$
621 $\varphi(T, S_0)$ for every Lebesgue measurable set B . Taking $B = \varphi(T, A) = \mathcal{R}_t(A)$ and using injectivity
622 of $x \mapsto \varphi(T, x)$ on $\overline{S_0}$, we have

$$\begin{aligned} |A| &= |\varphi(T, \cdot)^{-1}(\varphi(T, A))| \\ &\leq e^{nLT}|\mathcal{R}_T(A)|. \end{aligned}$$

623 Rearranging yields $|\mathcal{R}_T(A)| \geq e^{-nLT}|A|$, which proves (3). This completes the proof. \square

624 The preceding proposition is the technical tool needed to propagate the local thickness of the initial
625 set. We now combine its Lipschitz and volume-distortion bounds with the initial local Lebesgue-
626 mass bound from Proposition 1. This yields a corresponding local Lebesgue-mass lower bound for
627 the reachable set at time T .

628 **Proposition 2** (Propagation of local Lebesgue mass). Let $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$. Then, for every
629 $x \in \overline{S_T}$ and every $r \in (0, e^{LT}r_0]$, $|\mathcal{B}_r(x) \cap \overline{S_T}| \geq \omega_n 2^{-n} e^{-2nLT} r^n$.

630 *Proof of Proposition 2.* Fix $x \in \overline{S_T}$ and $r \in (0, e^{LT}r_0]$. Since $\overline{S_T} = \varphi(T, \overline{S_0})$ and the flow map
631 is injective on $\overline{S_0}$ by Proposition 4, there exists a unique $y \in \overline{S_0}$ such that $x = \varphi(T, y)$. Set
632 $\rho := e^{-LT}r$. Then $0 < \rho \leq r_0$. By Proposition 1,

$$|\mathcal{B}_\rho(y) \cap \overline{S_0}| \geq \omega_n 2^{-n} \rho^n.$$

633 Define

$$A := \mathcal{B}_\rho(y) \cap \overline{S_0}.$$

634 For every $p \in A$, the Lipschitz estimate in Proposition 4 gives

$$\|\varphi(T, p) - x\| = \|\varphi(T, p) - \varphi(T, y)\| \leq e^{LT}\|p - y\| \leq e^{LT}\rho = r.$$

635 Therefore, $\mathcal{R}_T(A) \subseteq \mathcal{B}_r(x) \cap \overline{S_T}$. Using the volume distortion lower bound in Proposition 4, we
636 obtain $|\varphi(T, A)| \geq e^{-nLT}|A|$. Hence,

$$\begin{aligned} |\mathcal{B}_r(x) \cap \overline{S_T}| &\geq |\varphi(T, A)| \\ &\geq e^{-nLT}|A| \\ &\geq e^{-nLT}\omega_n 2^{-n} \rho^n \\ &= \omega_n 2^{-n} e^{-2nLT} r^n, \end{aligned}$$

637 where the last equality uses $\rho = e^{-LT}r$. \square

638 The previous result propagates geometric thickness from $\overline{S_0}$ to $\overline{S_T}$. We also need to propagate the
639 sampling density lower bound. The next proposition shows that the pushforward distribution re-
640 mains absolutely continuous and that its density lower bound degrades by at most the same volume-
641 distortion factor.

642 **Proposition 3** (Propagation of density lower bounds). Suppose $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$ and Assump-
643 tion 1 holds. Then $P_T = (\varphi(T, \cdot))_{\#} P_0$ is absolutely continuous with respect to Lebesgue measure
644 on $\overline{S_T}$. Moreover, if p_T denotes its density, then

$$p_T(y) \geq \underline{p}_0 e^{-nLT}, \quad \text{for a.e. } y \in \overline{S_T}.$$

645 In particular,

$$\underline{p}_T := \underline{p}_0 e^{-nLT}$$

646 is a density lower bound for the endpoint distribution at time T .

647 *Proof of Proposition 3.* Let $B \subseteq \overline{S_T}$ be a Lebesgue measurable set and define

$$A := \varphi(T, \cdot)^{-1}(B) \subseteq \overline{S_0}.$$

648 Since P_T is the pushforward of P_0 under $\varphi(T, \cdot)$,

$$P_T(B) = P_0(A).$$

649 We first prove absolute continuity. By Proposition 4, the inverse map $\varphi(T, \cdot)^{-1} : \overline{S_T} \rightarrow \overline{S_0}$ is
650 e^{LT} -Lipschitz. Hence, for every Lebesgue measurable $B \subseteq \overline{S_T}$,

$$|A| = |\varphi(T, \cdot)^{-1}(B)| \leq e^{nLT}|B|.$$

651 Therefore, if $|B| = 0$, then $|A| = 0$. Since P_0 is absolutely continuous with respect to Lebesgue
652 measure on $\overline{S_0}$, this implies

$$P_T(B) = P_0(A) = 0.$$

653 Thus, P_T is absolutely continuous with respect to Lebesgue measure and admits a density p_T . It
654 remains to prove the density lower bound. Since the flow map is injective on $\overline{S_0}$, we have $B =$
655 $\mathcal{R}_T(A)$. By the upper Lipschitz bound in Proposition 4,

$$|B| = |\mathcal{R}_T(A)| \leq e^{nLT}|A|.$$

656 Therefore, $|A| \geq e^{-nLT}|B|$. Using the lower bound on the initial density, we get

$$P_T(B) = P_0(A) = \int_A p_0(x) dx \geq \underline{p}_0|A| \geq \underline{p}_0 e^{-nLT}|B|.$$

657 Since this holds for every Lebesgue measurable set $B \subseteq \overline{S_T}$, the density p_T satisfies

$$p_T(y) \geq \underline{p}_0 e^{-nLT}, \quad \text{for a.e. } y \in \overline{S_T}.$$

658 This proves the claim. □

659 We now combine the initial local Lebesgue-mass bound, the flow distortion estimate, and the end-
660 point density lower bound to prove the local probability lower bound used in the sample-complexity
661 argument.

662 **Corollary 2** (Endpoint local probability lower bound). Suppose $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$ and Assump-
663 tion 1 holds. Define

$$\Gamma_T(r) := \underline{p}_0 \omega_n 2^{-n} e^{-3nLT} r^n.$$

664 Then, for every $x \in \overline{S_T}$ and every $r \in (0, e^{LT}r_0]$,

$$P_T(\mathcal{B}_r(x) \cap \overline{S_T}) \geq \Gamma_T(r).$$

665 *Proof of Corollary 2.* By Proposition 3 and Proposition 3, we know for every $x \in \overline{S_T}$ and every
666 $r \in (0, e^{LT}r_0]$, $|\mathcal{B}_r(x) \cap \overline{S_T}| \geq \omega_n 2^{-n} e^{-2nLT} r^n$ and the endpoint density satisfies $p_T(z) \geq$
667 $\underline{p}_0 e^{-nLT}$, for a.e. $z \in \overline{S_T}$. Therefore,

$$\begin{aligned} P_T(\mathcal{B}_r(x) \cap \overline{S_T}) &= \int_{\mathcal{B}_r(x) \cap \overline{S_T}} p_T(z) dz \\ &\geq \underline{p}_0 e^{-nLT} |\mathcal{B}_r(x) \cap \overline{S_T}| \\ &\geq \underline{p}_0 e^{-nLT} \omega_n 2^{-n} e^{-2nLT} r^n \\ &= \underline{p}_0 \omega_n 2^{-n} e^{-3nLT} r^n \\ &= \Gamma_T(r). \end{aligned}$$

668 This proves the claim. □

669 **B.4 Sample Complexity**

670 **Sample Complexity Upper Bound.** We prove Theorem 1 using a direct random-covering argument,
 671 which only uses the fact that \widehat{S}_N contains all endpoint samples.

672 We first record the covering-number estimate used in the proof.

673 **Definition 3** (Covering number). *For a bounded set $A \subset \mathbb{R}^n$ and a radius $\rho > 0$, the ρ -covering*
 674 *number of A is defined as*

$$\mathcal{N}(A, \rho) := \min\{M \in \mathbb{N} : \exists z_1, \dots, z_M \in A \text{ such that } A \subseteq \bigcup_{j=1}^M \mathcal{B}_\rho(z_j)\}.$$

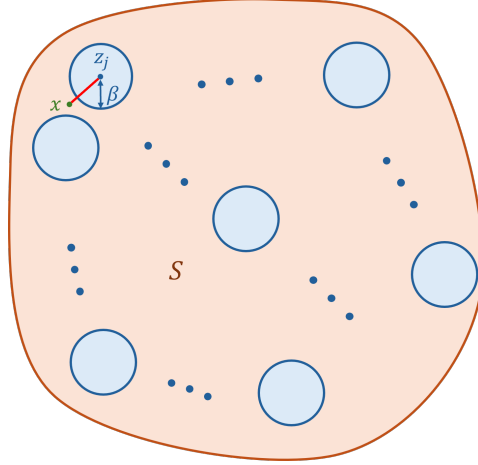


Figure 5: Set Covering

675 **Lemma 2** (Covering number of the reachable set). *Let $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$. For every $0 < \rho \leq$*
 676 *$e^{LT} r_0$,*

$$\mathcal{N}(\overline{S_T}, \rho) \leq \frac{2^{2n} e^{nLT}}{\omega_n \rho^n}.$$

677 *In particular, for $\rho = r/2$, $\mathcal{N}(\overline{S_T}, r/2) \leq \frac{2^{3n} e^{nLT}}{\omega_n r^n}$.*

678 *Proof.* Since the flow map $\varphi(T, \cdot)$ is e^{LT} -Lipschitz an injective on $\overline{S_0}$, any $e^{-LT} \rho$ -cover of $\overline{S_0}$ is
 679 mapped to a ρ -cover of $\overline{S_T}$. Hence

$$\mathcal{N}(\overline{S_T}, \rho) \leq \mathcal{N}(\overline{S_0}, e^{-LT} \rho).$$

680 It remains to bound the covering number of $\overline{S_0}$. Fix $\beta \in (0, r_0]$. Let $\{z_1, \dots, z_M\} \subset \overline{S_0}$ be a
 681 maximal β -separated set, meaning that $\|z_i - z_j\| > \beta, i \neq j$, and no additional point of $\overline{S_0}$ can be
 682 added while preserving this property. By maximality, this set is a β -net of $\overline{S_0}$. Indeed, as Figure 5
 683 shows, if there existed $x \in \overline{S_0}$ such that

$$x \notin \bigcup_{j=1}^M \mathcal{B}_\beta(z_j),$$

684 then $\|x - z_j\| > \beta$ for every j , so $\{z_1, \dots, z_M, x\}$ would still be β -separated, contradicting maxi-
 685 mality. Therefore, $\overline{S_0} \subseteq \bigcup_{j=1}^M \mathcal{B}_\beta(z_j)$, and hence

$$\mathcal{N}(\overline{S_0}, \beta) \leq M.$$

686 We now upper bound M . Since the points z_1, \dots, z_M are β -separated, the balls $\mathcal{B}_{\beta/2}(z_j)$ are pair-
 687 wise disjoint up to their boundaries. Therefore, the sets

$$\mathcal{B}_{\beta/2}(z_j) \cap S_0, \quad j = 1, \dots, M,$$

688 are also pairwise disjoint. By Proposition 1, applied with radius $\beta/2$, we have

$$|\mathcal{B}_{\beta/2}(z_j) \cap S_0| \geq \omega_n 2^{-n} \left(\frac{\beta}{2}\right)^n = \omega_n 2^{-2n} \beta^n, \quad j = 1, \dots, M.$$

689 Since these sets are pairwise disjoint subsets of S_0 , and $|S_0| = 1$ in \mathcal{F}_{L, r_0} ,

$$1 = |S_0| \geq \sum_{j=1}^M |\mathcal{B}_{\beta/2}(z_j) \cap S_0| \geq M \omega_n 2^{-2n} \beta^n.$$

690 Thus, $M \leq \frac{2^{2n}}{\omega_n \beta^n}$. Combining this with $\mathcal{N}(\overline{S_0}, \beta) \leq M$ gives

$$\mathcal{N}(\overline{S_0}, \beta) \leq \frac{2^{2n}}{\omega_n \beta^n}, \quad 0 < \beta \leq r_0.$$

691 Finally, take $\beta = e^{-LT} \rho$. Since $\rho \leq e^{LT} r_0$, we have $\beta \leq r_0$, and therefore

$$\mathcal{N}(\overline{S_T}, \rho) \leq \mathcal{N}(\overline{S_0}, e^{-LT} \rho) \leq \frac{2^{2n}}{\omega_n (e^{-LT} \rho)^n} = \frac{2^{2n} e^{nLT}}{\omega_n \rho^n}.$$

692 Setting $\rho = r/2$ yields

$$\mathcal{N}(\overline{S_T}, r/2) \leq \frac{2^{3n} e^{nLT}}{\omega_n r^n}.$$

693

□

694 We now prove the sampling upper bound. The proof combines the endpoint local probability lower
695 bound from Corollary 2 with a covering argument: we cover $\overline{S_T}$ by an $r/2$ -net and show that, with
696 high probability, every net ball is hit by at least one endpoint sample.

697 **Theorem 1** (Sampling upper bound for inner coverage). Let $(S_0, F, P_0) \in \mathcal{F}_{L, r_0}$, and let
698 $Y_1, \dots, Y_N \stackrel{\text{i.i.d.}}{\sim} P_T$. Suppose that Assumption 1 holds and that the estimator \widehat{S}_N contains all
699 endpoint samples, i.e., $Y_i \in \widehat{S}_N$ for all $i = 1, \dots, N$. Fix $0 < r \leq 2e^{LT} r_0$. If

$$N \geq \frac{2^{2n} e^{3nLT}}{\underline{p}_0 \omega_n r^n} \left[\log \left(\frac{2^{3n} e^{nLT}}{\omega_n r^n} \right) + \log \frac{1}{\delta} \right],$$

700 then, with probability at least $1 - \delta$,

$$\sup_{x \in \overline{S_T}} d(x, \widehat{S}_N) \leq r.$$

701 *Proof of Theorem 1.* Let

$$q_r := \Gamma_T(r/2) = \underline{p}_0 \omega_n 2^{-n} e^{-3nLT} \left(\frac{r}{2}\right)^n = \underline{p}_0 \omega_n 2^{-2n} e^{-3nLT} r^n$$

702 and $\{z_1, \dots, z_M\} \subset \overline{S_T}$ be an $r/2$ -net of $\overline{S_T}$. By Lemma 2, we may choose this net so that
703 $M \leq \frac{2^{3n} e^{nLT}}{\omega_n r^n}$. By Corollary 2, for each $j = 1, \dots, M$,

$$P_T(\mathcal{B}_{r/2}(z_j) \cap \overline{S_T}) \geq \Gamma_T(r/2) = q_r.$$

704 Since P_T is supported on $\overline{S_T}$, this also gives $P_T(\mathcal{B}_{r/2}(z_j)) \geq q_r$. For each j , define the bad event

$$E_j := \{\mathcal{B}_{r/2}(z_j) \cap \{Y_1, \dots, Y_N\} = \emptyset\}.$$

705 Let $\mathbb{P}(E_j)$ represents the probability that E_j happens. Because Y_1, \dots, Y_N are independent samples
706 from P_T ,

$$\mathbb{P}(E_j) = (1 - P_T(\mathcal{B}_{r/2}(z_j)))^N \leq (1 - q_r)^N \leq e^{-Nq_r}.$$

707 Therefore, by the union bound,

$$\mathbb{P}\left(\bigcup_{j=1}^M E_j\right) \leq \sum_{j=1}^M \mathbb{P}(E_j) \leq M e^{-Nq_r}.$$

708 Then suppose sample-size $N \geq \frac{2^{2n} e^{3nLT}}{\underline{p}_0 \omega_n r^n} \left[\log \left(\frac{2^{3n} e^{nLT}}{\omega_n r^n} \right) + \log \frac{1}{\delta} \right]$, because $\frac{1}{q_r} = \frac{2^{2n} e^{3nLT}}{\underline{p}_0 \omega_n r^n}$ and
709 $M \leq \frac{2^{3n} e^{nLT}}{\omega_n r^n}$. Thus, $M e^{-Nq_r} \leq \delta$. Hence, with probability at least $1 - \delta$, none of the events E_j
710 occurs. Equivalently, every $r/2$ -net ball contains at least one endpoint sample. On this event, fix
711 any $x \in \overline{S_T}$. Since $\{z_1, \dots, z_M\}$ is an $r/2$ -net, there exists z_j such that $\|x - z_j\| \leq r/2$. Since
712 E_j does not occur, there exists a sample Y_i such that $\|Y_i - z_j\| \leq r/2$. By the sample-containment
713 assumption, $Y_i \in \widehat{S}_N$. Therefore,

$$d(x, \widehat{S}_N) \leq \|x - Y_i\| \leq \|x - z_j\| + \|z_j - Y_i\| \leq r.$$

714 Since this holds for every $x \in \overline{S_T}$, we conclude that

$$\sup_{x \in \overline{S_T}} d(x, \widehat{S}_N) = \sup_{x \in \overline{S_T}} d(x, \widehat{S}_N) \leq r.$$

715 This completes the proof. □

716 The preceding theorem controls only the inner directed error from the reachable set to the estimator.
717 We now add the estimator-dependent outer deviation to obtain a full Hausdorff bound.

718 **Corollary 3** (Hausdorff bound with outer deviation). Under the conditions of Theorem 1, suppose
719 in addition that the estimator satisfies the outer-deviation condition

$$\widehat{S}_N \subseteq \mathcal{B}_\eta(S_T)$$

720 for some $\eta > 0$. Then, with probability at least $1 - \delta$,

$$d_H(S_T, \widehat{S}_N) \leq \max\{r, \eta\}.$$

721 *Proof of Corollary 3.* By Theorem 1, with probability at least $1 - \delta$, $\sup_{x \in \overline{S_T}} d(x, \widehat{S}_N) \leq r$. The
722 outer-deviation condition $\widehat{S}_N \subseteq \mathcal{B}_\eta(S_T)$ implies $\sup_{z \in \widehat{S}_N} d(z, S_T) \leq \eta$. Therefore, on the same
723 event,

$$\begin{aligned} d_H(S_T, \widehat{S}_N) &= \max\left\{ \sup_{x \in S_T} d(x, \widehat{S}_N), \sup_{z \in \widehat{S}_N} d(z, S_T) \right\} \\ &\leq \max\{r, \eta\}. \end{aligned}$$

724 Thus, with probability at least $1 - \delta$,

$$d_H(S_T, \widehat{S}_N) \leq \max\{r, \eta\}.$$

725 □

726 **Sample Complexity Lower Bound.** To obtain the sample-complexity lower bound, we first identify
727 the local statistical obstruction that any sampling-based method must face. The upper bound relies
728 on the fact that every r -scale neighborhood of the reachable set has nontrivial probability-mass. For
729 the lower bound, we need the converse type of phenomenon: among many disjoint r -scale regions
730 inside a regular support, at least one region must have small probability-mass. Otherwise, the total
731 probability over all such regions would exceed one.

732 We formalize this idea through an interior packing number.

733 **Definition 4** (Interior packing number). Let $S_0 \subset \mathbb{R}^n$ be a nonempty set and let $r > 0$. The interior
734 r -packing number of S_0 , denoted by $M_r(S_0)$, is the largest integer M for which there exist points
735 $z_1, \dots, z_M \in S_0$ such that

$$\mathcal{B}_r(z_j) \subset S_0, \quad \mathcal{B}_r(z_i) \cap \mathcal{B}_r(z_j) = \emptyset \quad (i \neq j).$$

736 Equivalently, $M_r(S_0)$ is the maximum number of pairwise disjoint closed balls of radius r that can
737 be placed entirely inside S_0 .

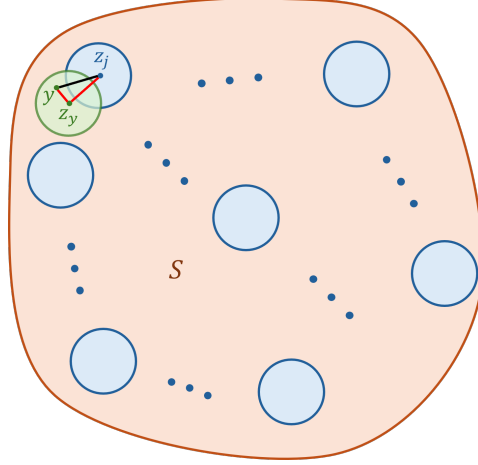


Figure 6: Interior Packing

738 Positive reach provides the geometric ingredient needed to lower bound this packing number. It
 739 ensures that the support has enough interior thickness to contain order r^{-n} disjoint balls of radius r .
 740 A simple averaging argument then shows that, for any probability measure supported on S_0 , at least
 741 one of these balls has probability-mass at most order r^n .

742 **Lemma 3** (Interior packing from positive reach). *Let $S_0 \subset \mathbb{R}^n$ be a bounded open set satisfying*
 743 *$\text{reach}(S_0^c) \geq r_0 > 0$ and $|S_0| = 1$. Then, for every $0 < r \leq r_0$, there exist points $z_1, \dots, z_{M_r} \in S_0$*
 744 *such that $\mathcal{B}_r(z_j) \subset S_0$, $\mathcal{B}_r(z_i) \cap \mathcal{B}_r(z_j) = \emptyset$ ($i \neq j$), and $M_r \geq \frac{3^{-n}r^{-n}}{\omega_n}$.*

745 *Proof.* Since $\text{reach}(S_0^c) \geq r_0$, the opening property implies that $S_0 = S_0 \circ \mathcal{B}_r(0)$ for every $0 <$
 746 $r \leq r_0$. Hence, as Figure 6 shows, for any $r \in (0, r_0)$ and for every $y \in S_0$, there exists a point
 747 $z_y \in S_0 \ominus \mathcal{B}_r(0)$ such that $y \in \mathcal{B}_r(z_y)$. In particular,

$$\mathcal{B}_r(z_y) \subset S_0.$$

748 Let $\{\mathcal{B}_r(z_j)\}_{j=1}^{M_r}$ be a maximal family of pairwise disjoint balls contained in S_0 . By maximality,
 749 for every $y \in S_0$, the interior ball $\mathcal{B}_r(z_y) \subset S_0$ must intersect at least one selected ball $\mathcal{B}_r(z_j)$;
 750 otherwise it could be added to the family. Thus $\|z_y - z_j\| \leq 2r$ for some j . Since $y \in \mathcal{B}_r(z_y)$, we
 751 get

$$\|y - z_j\| \leq \|y - z_y\| + \|z_y - z_j\| \leq 3r.$$

752 Therefore,

$$S_0 \subseteq \bigcup_{j=1}^{M_r} \mathcal{B}_{3r}(z_j).$$

753 Taking Lebesgue measure gives

$$1 = |S_0| \leq \sum_{j=1}^{M_r} |\mathcal{B}_{3r}(z_j)| = M_r \omega_n (3r)^n.$$

754 Rearranging yields

$$M_r \geq \frac{3^{-n}r^{-n}}{\omega_n}.$$

755 This proves the claim. \square

756 The lemma above is purely geometric: it lower bounds the number of disjoint r -scale balls that can
 757 be placed inside a positive-reach set. We now turn this packing statement into a probabilistic one.
 758 Since the balls are disjoint and the total probability-mass is one, at least one of these balls must carry
 759 probability no larger than the average mass $1/M_r$.

760 **Corollary 4** (Existence of a low-probability ball). *Under the conditions of Lemma 3, for any prob-*
 761 *ability measure P_0 supported on $\overline{S_0}$, there exists $j \in \{1, \dots, M_r\}$ such that*

$$P_0(\mathcal{B}_r(z_j)) \leq \frac{1}{M_r} \leq 3^n \omega_n r^n.$$

762 *Proof.* Since the balls $\mathcal{B}_r(z_1), \dots, \mathcal{B}_r(z_{M_r})$ are pairwise disjoint, $\bigcup_{j=1}^{M_r} \mathcal{B}_r(z_j) \subseteq \overline{S_0}$ and P_0 is a
 763 probability measure on the Borel σ -algebra of \mathbb{R}^n supported on $\overline{S_0}$,

$$1 = P_0(\overline{S_0}) \geq P_0\left(\bigcup_{j=1}^{M_r} \mathcal{B}_r(z_j)\right) = \sum_{j=1}^{M_r} P_0(\mathcal{B}_r(z_j)).$$

764 Hence at least one term satisfies $P_0(\mathcal{B}_r(z_j)) \leq 1/M_r$. Using Lemma 3 gives the fact that there
 765 exists $j \in \{1, \dots, M_r\}$ such that

$$P_0(\mathcal{B}_r(z_j)) \leq \frac{1}{M_r} \leq 3^n \omega_n r^n.$$

766

□

767 This low-probability ball will be removed in the two-point construction below, producing a
 768 Hausdorff-visible perturbation that is statistically difficult to detect. More precisely, we will con-
 769 struct two admissible reachability instances whose endpoint supports are close in distribution but far
 770 in Hausdorff distance. The two instances differ only through the initial domain and sampling law:
 771 the first uses the original initial domain, while the second removes the low-probability interior ball
 772 from it. After propagation through the flow, these two initial domains induce two endpoint supports.

773 To make the testing reduction independent of this particular construction, we first state it for two
 774 generic candidate endpoint supports.

775 **Definition 5** (Hypothesis test induced by a support estimator). *Let $S^{(0)}, S^{(1)} \subset \mathbb{R}^n$ be two can-*
 776 *didate endpoint supports. Given a support estimator \widehat{S}_N , define the induced binary test $\psi =$
 777 $\psi(Y_1, \dots, Y_N) \in \{0, 1\}$ by*

$$\psi(Y_1, \dots, Y_N) := \begin{cases} 0, & d_H(\widehat{S}_N, S^{(0)}) \leq d_H(\widehat{S}_N, S^{(1)}), \\ 1, & d_H(\widehat{S}_N, S^{(1)}) < d_H(\widehat{S}_N, S^{(0)}). \end{cases}$$

778 The test simply chooses the candidate endpoint support that is closer to the estimated set in Haus-
 779 dorff distance. Therefore, if the estimator is accurate and the two candidate endpoint supports are
 780 separated by more than twice the target accuracy, the induced test must also be correct. This obser-
 781 vation is formalized next.

782 **Lemma 4** (Accurate support estimation induces a correct test). *Suppose $d_H(S^{(0)}, S^{(1)}) > 2r$, and*
 783 *let ψ be the test in Definition 5. If the true endpoint support is $S^{(0)}$ and*

$$d_H(\widehat{S}_N, S^{(0)}) \leq r,$$

784 *then $\psi = 0$. Similarly, if the true endpoint support is $S^{(1)}$ and*

$$d_H(\widehat{S}_N, S^{(1)}) \leq r,$$

785 *then $\psi = 1$.*

786 *Proof.* If $d_H(\widehat{S}_N, S^{(0)}) \leq r$, then by the triangle inequality,

$$d_H(\widehat{S}_N, S^{(1)}) \geq d_H(S^{(0)}, S^{(1)}) - d_H(\widehat{S}_N, S^{(0)}) > 2r - r = r.$$

787 Hence $d_H(\widehat{S}_N, S^{(0)}) < d_H(\widehat{S}_N, S^{(1)})$, and therefore $\psi = 0$ by Definition 5. The argument for $S^{(1)}$
 788 is identical. □

789 To compare two such reachable-set instances statistically, we need to quantify how distinguishable
790 the corresponding endpoint-sampling distributions are. This motivates the following two standard
791 information-theoretic notions. The Radon–Nikodym derivative allows us to express one probability
792 measure as a density with respect to another, and the Kullback–Leibler divergence then measures
793 the information gap between the two induced sampling models.

794 **Definition 6** (Radon–Nikodym derivative [58]). *Let P and Q be probability measures on a measur-*
795 *able space (Ω, \mathcal{F}) . We say that P is absolutely continuous with respect to Q , denoted by $P \ll Q$, if*
796 *$Q(A) = 0$ implies $P(A) = 0$ for every $A \in \mathcal{F}$. In this case, the Radon–Nikodym derivative $\frac{dP}{dQ}$ is*
797 *the measurable function satisfying*

$$P(A) = \int_A \frac{dP}{dQ} dQ, \quad \forall A \in \mathcal{F}.$$

798 **Definition 7** (Kullback–Leibler divergence [59]). *Let P and Q be probability measures on (Ω, \mathcal{F}) .*
799 *If $P \ll Q$, the Kullback–Leibler divergence from P to Q is*

$$D_{\text{KL}}(P\|Q) := \int_{\Omega} \log \left(\frac{dP}{dQ} \right) dP.$$

800 *If $P \not\ll Q$, we set $D_{\text{KL}}(P\|Q) = +\infty$.*

801 The preceding lemma reduces Hausdorff estimation to binary testing. We now use a standard testing
802 lower bound to show that this binary decision problem cannot be solved reliably when the two
803 sampling distributions have small KL divergence.

804 **Lemma 5** ([60], Lemma 2.6). *Let P and Q be probability measures with $P \ll Q$. Then, for any*
805 *binary test $\psi \in \{0, 1\}$,*

$$P(\psi = 1) + Q(\psi = 0) \geq \frac{1}{2} \exp(-D_{\text{KL}}(P\|Q)).$$

806 *Consequently, for N i.i.d. samples,*

$$P^N(\psi = 1) + Q^N(\psi = 0) \geq \frac{1}{2} \exp(-D_{\text{KL}}(P^N\|Q^N)).$$

807 With these ingredients in place, we now prove the minimax lower bound. The proof constructs two
808 admissible instances

$$(S_0^{(0)}, F^{(0)}, P_0^{(0)}) \quad \text{and} \quad (S_0^{(1)}, F^{(1)}, P_0^{(1)})$$

809 in \mathcal{F}_{L,r_0} . The second instance is obtained by removing a low-probability interior ball from the initial
810 domain of the first. Under an admissible expanding flow, the resulting endpoint supports $\overline{S}_T^{(0)}$ and
811 $\overline{S}_T^{(1)}$ are separated in Hausdorff distance, while the endpoint sampling distributions $P_T^{(0)}$ and $P_T^{(1)}$
812 have small KL divergence.

813 We now prove the minimax lower bound. The proof uses a two-point testing argument. We construct
814 two admissible instances in the problem family whose endpoint supports are separated in Hausdorff
815 distance, but whose endpoint sampling distributions are close in KL divergence. Any estimator that
816 is Hausdorff-accurate on both instances would therefore induce a reliable binary test, contradicting
817 the testing lower bound when the sample size is too small.

818 **Theorem 2** (Minimax Hausdorff lower bound). *Let $C := 3^n \omega_n$. For any accuracy $0 < r \leq$*
819 *$2^{-\frac{n+1}{n}} e^{LT} r_0$ and any confidence level $1 - \delta$ with $\delta \in (0, 1)$, if*

$$N < \frac{e^{nLT}}{2^{n+1} C r^n} \log \frac{1}{4\delta}, \quad \text{and} \quad r_0 \leq C^{-\frac{1}{n}}$$

820 *then for every estimator \widehat{S}_N , there exists an instance $(S_0, F, P_0) \in \mathcal{F}_{L,r_0}$ such that*

$$P_T^N \left(d_H(S_T, \widehat{S}_N) > r \right) > \delta.$$

821 *Proof of Theorem 2.* We prove the result by a two-point minimax testing argument. Fix an arbitrary
822 estimator \hat{S}_N . It suffices to construct two admissible L -Lipschitz dynamics and initial distributions
823 whose endpoint supports are separated in Hausdorff distance, but whose endpoint sample distribu-
824 tions are statistically close. To be more specific, our goal is to construct two admissible instances
825 $(S_0^{(0)}, F, P_0^{(0)})$ and $(S_0^{(1)}, F, P_0^{(1)})$ are in \mathcal{F}_{L, r_0} whose endpoint supports, denoted by $\overline{S_T^{(0)}}$ and
826 $\overline{S_T^{(1)}}$, are separated in Hausdorff distance, while their endpoint sampling distributions $P_T^{(0)}$ and
827 $P_T^{(1)}$ are statistically close.

828 Let $R := \omega_n^{-1/n}$ and choose $x_0 \in \mathbb{R}^n$. Define the reference initial domain by $S_0^{(0)} :=$
829 $\mathcal{B}_R^\circ(x_0), \overline{S_0^{(0)}} := \overline{S_0^{(0)}} = \mathcal{B}_R(x_0), P_0^{(0)} := \text{Unif}(S_0^{(0)})$. Then $|S_0^{(0)}| = 1$ and $\text{supp}(P_0^{(0)}) = \overline{S_0^{(0)}}$.
830 Moreover, since $R = \omega_n^{-1/n}$ and $r_0 \leq C^{-1/n} = \frac{1}{3}\omega_n^{-1/n}$, we have $\text{reach}((S_0^{(0)})^c) \geq r_0$ and
831 $S_0^{(0)} \ominus \mathcal{B}_{3r_0}(0) \neq \emptyset$.

832 We use the admissible L -Lipschitz vector field $F(x) = Lx$, whose flow is $\varphi(T, x) = e^{LT}x$. In
833 particular, $\varphi(T, \mathcal{B}_\rho(z)) = \mathcal{B}_{e^{LT}\rho}(\varphi(T, z)), \forall \rho > 0, z \in \mathbb{R}^n$. Fix $\alpha > 0$ sufficiently small and
834 set $\rho := (2 + \alpha)e^{-LT}r$. By the assumed upper range on r , we may choose α sufficiently small
835 so that $\rho = (2 + \alpha)e^{-LT}r \leq r_0$. Applying Corollary 4 to $S_0^{(0)}$ at radius ρ , we obtain a point
836 $z_j \in S_0^{(0)} \ominus \mathcal{B}_{3r_0}(0)$ such that $\mathcal{B}_\rho(z_j) \subset S_0^{(0)}$ and

$$q_r := P_0^{(0)}(\mathcal{B}_\rho^\circ(z_j)) \leq P_0^{(0)}(\mathcal{B}_\rho(z_j)) \leq C\rho^n = (2 + \alpha)^n C e^{-nLT} r^n.$$

837 The condition $r \leq 2^{-(n+1)/n} e^{LT} r_0$ together with $r_0 \leq C^{-1/n}$ implies $q_r \leq \frac{1}{2}(\frac{2+\alpha}{2})^n$.

838 The perturbed initial domain is defined by removing this closed ball: $S_0^{(1)} := S_0^{(0)} \setminus \mathcal{B}_\rho(z_j)$. Since
839 $S_0^{(0)}$ is open and $\mathcal{B}_\rho(z_j)$ is closed, $S_0^{(1)}$ is open. Its corresponding support is $\overline{S_0^{(1)}} := \overline{S_0^{(1)}} =$
840 $\overline{S_0^{(0)}} \setminus \mathcal{B}_\rho^\circ(z_j)$. We define the second initial distribution by conditioning on this new support: $P_0^{(1)} :=$
841 $P_0^{(0)}(\cdot | \overline{S_0^{(1)}})$. Equivalently, up to boundary sets of Lebesgue measure zero, $P_0^{(1)} = \text{Unif}(S_0^{(1)})$,
842 and hence $\text{supp}(P_0^{(1)}) = \overline{S_0^{(1)}}$. It remains to verify the reach condition for the perturbed domain.
843 Since

$$(S_0^{(1)})^c = (S_0^{(0)})^c \cup \mathcal{B}_\rho(z_j),$$

844 and $z_j \in S_0^{(0)} \ominus \mathcal{B}_{3r_0}(0)$ while $\rho \leq r_0$, the two closed components $(S_0^{(0)})^c$ and $\mathcal{B}_\rho(z_j)$ are separated
845 by distance at least $2r_0$. Moreover, $(S_0^{(0)})^c$ has reach at least r_0 , and $\mathcal{B}_\rho(z_j)$ is convex. Therefore,
846 by the separated-union reach lemma,

$$\text{reach}((S_0^{(1)})^c) \geq r_0.$$

847 Thus both $(S_0^{(0)}, F, P_0^{(0)})$ and $(S_0^{(1)}, F, P_0^{(1)})$ belong to \mathcal{F}_{L, r_0} .

848 By the L -exponential growth property, $\varphi(T, \mathcal{B}_{2e^{-LT}r}(z_j)) = \mathcal{B}_{2r}(y_j)$. Thus a low-probability ini-
849 tial ball at scale $2e^{-LT}r$ becomes an endpoint ball at Hausdorff scale $2r$. We first verify the Haus-
850 dorff separation between the two endpoint supports. Recall that $\overline{S_T^{(0)}} := \varphi(T, \overline{S_0^{(0)}}), \overline{S_T^{(1)}} :=$
851 $\varphi(T, \overline{S_0^{(1)}})$, where $\overline{S_0^{(1)}} = \overline{S_0^{(0)}} \setminus \mathcal{B}_\rho^\circ(z_j)$. Since the flow is $\varphi(T, x) = e^{LT}x$ and $\rho =$
852 $(2 + \alpha)e^{-LT}r$, we have $\overline{S_T^{(1)}} = \overline{S_T^{(0)}} \setminus \mathcal{B}_{2r}^\circ(y_j), y_j := \varphi(T, z_j)$. In particular, $y_j \in \overline{S_T^{(0)}}$
853 and the closest points of $\overline{S_T^{(1)}}$ to y_j lie on the boundary of $\mathcal{B}_{2r}(y_j)$. Hence

$$d(y_j, \overline{S_T^{(1)}}) = (2 + \alpha)r > 2r. \quad (4)$$

854 Therefore, $d_H(\overline{S_T^{(0)}}, \overline{S_T^{(1)}}) \geq d(y_j, \overline{S_T^{(1)}}) > 2r$. Next, we bound the KL divergence between the
855 two endpoint distributions. Since

$$P_0^{(1)} = P_0^{(0)}(\cdot | \overline{S_0^{(1)}}), \quad P_0^{(0)}(\overline{S_0^{(1)}}) = 1 - q_r,$$

856 we have, for every measurable set $A \subseteq \mathbb{R}^n$,

$$P_0^{(1)}(A) = \frac{P_0^{(0)}(A \cap \overline{S_0^{(1)}})}{1 - q_r}.$$

857 Equivalently, this conditional distribution can be written in Radon–Nikodym form as $P_0^{(1)}(A) =$
 858 $\int_A \frac{\mathbf{1}_{\overline{S_0}^{(1)}}(x)}{1 - q_r} dP_0^{(0)}(x)$. Therefore,

$$\frac{dP_0^{(1)}}{dP_0^{(0)}}(x) = \frac{\mathbf{1}_{\overline{S_0}^{(1)}}(x)}{1 - q_r}, \quad P_0^{(0)}\text{-a.s.}$$

859 We now compute the KL divergence. By definition,

$$D_{\text{KL}}(P_0^{(1)} \| P_0^{(0)}) = \int \log \left(\frac{dP_0^{(1)}}{dP_0^{(0)}}(x) \right) dP_0^{(1)}(x).$$

860 Since $P_0^{(1)}$ is supported on $\overline{S_0}^{(1)}$, we have $\mathbf{1}_{\overline{S_0}^{(1)}}(x) = 1$ for $P_0^{(1)}$ -almost every x . Hence, on the
 861 support of $P_0^{(1)}$,

$$\frac{dP_0^{(1)}}{dP_0^{(0)}}(x) = \frac{1}{1 - q_r}.$$

862 Therefore,

$$\begin{aligned} D_{\text{KL}}(P_0^{(1)} \| P_0^{(0)}) &= \int_{\overline{S_0}^{(1)}} \log \left(\frac{1}{1 - q_r} \right) dP_0^{(1)}(x) \\ &= \log \left(\frac{1}{1 - q_r} \right) \int_{\overline{S_0}^{(1)}} dP_0^{(1)}(x) \\ &= \log \left(\frac{1}{1 - q_r} \right) P_0^{(1)}(\overline{S_0}^{(1)}) \\ &= \log \left(\frac{1}{1 - q_r} \right) \\ &= -\log(1 - q_r). \end{aligned} \tag{5}$$

863 Because $\varphi(T, \cdot)$ is a one-to-one flow map on the supports under consideration in the hard instance
 864 (e.g. $\dot{x} = Lx$), KL divergence is invariant under pushforward by $\varphi(T, \cdot)$. Thus

$$\begin{aligned} D_{\text{KL}}(P_T^{(1)} \| P_T^{(0)}) &= -\log(1 - q_r) \\ &\leq 2q_r \leq 2(2 + \alpha)^n C e^{-nLT} r^n, \end{aligned}$$

865 where we use (5) for the equation and the inequality holds when $q_r \leq \frac{1}{2} \left(\frac{2+\alpha}{2} \right)^n \leq 0.79$ for suffi-
 866 ciently small α . For N independent endpoint samples,

$$D_{\text{KL}} \left((P_T^{(1)})^N \| (P_T^{(0)})^N \right) = N D_{\text{KL}}(P_T^{(1)} \| P_T^{(0)}) \leq 2(2 + \alpha)^n N C e^{-nLT} r^n. \tag{6}$$

867 We now relate Hausdorff estimation to binary testing. Define the test induced by \widehat{S}_N as

$$\psi(Y_1, \dots, Y_N) := \begin{cases} 0, & d_H(\widehat{S}_N, \overline{S_T}^{(0)}) \leq d_H(\widehat{S}_N, \overline{S_T}^{(1)}), \\ 1, & d_H(\widehat{S}_N, \overline{S_T}^{(1)}) < d_H(\widehat{S}_N, \overline{S_T}^{(0)}). \end{cases}$$

868 By (4), if $d_H(\widehat{S}_N, \overline{S_T}^{(0)}) \leq r$, then $\psi = 0$; similarly, if $d_H(\widehat{S}_N, \overline{S_T}^{(1)}) \leq r$, then $\psi = 1$. In other
 869 words, the following two implications hold:

$$\begin{cases} d_H(\widehat{S}_N, \overline{S_T}^{(0)}) \leq r \implies \psi = 0, \\ d_H(\widehat{S}_N, \overline{S_T}^{(1)}) \leq r \implies \psi = 1. \end{cases}$$

870 Therefore, their contrapositives also hold. Since $\psi \in \{0, 1\}$, these contrapositives can be written
 871 explicitly as

$$\begin{cases} \psi = 1 \implies d_H(\widehat{S}_N, \overline{S_T}^{(0)}) > r, \\ \psi = 0 \implies d_H(\widehat{S}_N, \overline{S_T}^{(1)}) > r. \end{cases}$$

872 Equivalently, this gives the event inclusions

$$\begin{cases} \{\psi = 1\} \subseteq \{d_H(\widehat{S}_N, \overline{S}_T^{(0)}) > r\}, \\ \{\psi = 0\} \subseteq \{d_H(\widehat{S}_N, \overline{S}_T^{(1)}) > r\}. \end{cases}$$

873 Consequently,

$$(P_T^{(0)})^N \left(d_H(\widehat{S}_N, \overline{S}_T^{(0)}) > r \right) \geq (P_T^{(0)})^N(\psi = 1),$$

874 and

$$(P_T^{(1)})^N \left(d_H(\widehat{S}_N, \overline{S}_T^{(1)}) > r \right) \geq (P_T^{(1)})^N(\psi = 0).$$

875 Therefore, for this estimator \widehat{S}_N ,

$$\begin{aligned} \sup_{\ell \in \{0,1\}} (P_T^{(\ell)})^N \left(d_H(\widehat{S}_N, \overline{S}_T^{(\ell)}) > r \right) &\geq \max \left\{ (P_T^{(0)})^N(\psi = 1), (P_T^{(1)})^N(\psi = 0) \right\} \\ &\geq \frac{1}{2} \left[(P_T^{(0)})^N(\psi = 1) + (P_T^{(1)})^N(\psi = 0) \right]. \end{aligned} \quad (7)$$

876 By Lemma 5,

$$(P_T^{(0)})^N(\psi = 1) + (P_T^{(1)})^N(\psi = 0) \geq \frac{1}{2} \exp\left(-D_{\text{KL}}\left((P_T^{(1)})^N \parallel (P_T^{(0)})^N\right)\right).$$

877 Combining this with (7) and (6) gives

$$\sup_{\ell \in \{0,1\}} (P_T^{(\ell)})^N \left(d_H(\widehat{S}_N, \overline{S}_T^{(\ell)}) > r \right) \geq \frac{1}{4} \exp(-2(2 + \alpha)^n N C e^{-nLT} r^n). \quad (8)$$

878 Now assume

$$N < \frac{e^{nLT}}{2(2 + \alpha)^n C r^n} \log \frac{1}{4\delta}.$$

879 Then

$$2(2 + \alpha)^n N C e^{-nLT} r^n < \log \frac{1}{4\delta}.$$

880 Substituting this into (8) yields

$$\sup_{\ell \in \{0,1\}} (P_T^{(\ell)})^N \left(d_H(\widehat{S}_N, \overline{S}_T^{(\ell)}) > r \right) > \delta.$$

881 Since \widehat{S}_N was arbitrary, this means that for every estimator \widehat{S}_N , at least one of the two admissible
882 instances $(S_0^{(i)}, F, P_0^{(i)})$, $i = 1, 2$ satisfies

$$P_T^N \left(d_H(\widehat{S}_N, S_T) > r \right) = P_T^N \left(d_H(\widehat{S}_N, \overline{S}_T) > r \right) > \delta.$$

883 And such conclusion holds for all sufficiently small α , then we have when

$$N < \lim_{\alpha \rightarrow 0^+} \frac{e^{nLT}}{2(2 + \alpha)^n C r^n} \log \frac{1}{4\delta} = \frac{e^{nLT}}{2^{n+1} C r^n} \log \frac{1}{4\delta}.$$

884 There always exists an admissible L -Lipschitz dynamic class $F(\cdot)$, with endpoint support \overline{S}_T and
885 endpoint distribution P_T , such that $P_T^N \left(d_H(\widehat{S}_N, \overline{S}_T) > r \right) > \delta$, which proves the minimax lower
886 bound. \square

887 **Remark 2** (Worst-case and instance-dependent time dependence). *A natural question is whether the*
888 *exponential-in-time factor in Theorem 2 is only a pathological worst-case effect. The minimax result*
889 *is a uniform impossibility statement: it does not imply that every fixed instance requires sample*
890 *complexity scaling as $e^{nLT} r^{-n}$, where L is a global Lipschitz constant and T is the fixed time*

891 horizon. For a fixed instance, the relevant quantity is the smallest endpoint mass of an r -scale
 892 neighborhood,

$$\gamma_T(r) := \inf_{y \in \overline{S_T}} P_T(\mathcal{B}_r(y) \cap \overline{S_T}),$$

893 and the effective sample complexity is governed by $1/\gamma_T(r)$.

894 The lower-bound construction shows that exponential degradation can occur when the dynamics
 895 expands small initial neighborhoods before time T . In the hard instance used in the proof, an
 896 endpoint ball of radius r has a preimage with radius on the order of $e^{-LT}r$. Consequently, its initial
 897 probability mass is on the order of $e^{-nLT}r^n$, which leads to the lower-bound scaling $N \gtrsim e^{nLT}r^{-n}$
 898 up to constants.

899 For a general nonlinear system $\dot{x} = F(x)$, the corresponding instance-dependent quantity is con-
 900 trolled by the derivative of the flow map $D_x\varphi(T, x)$. If the flow expands volume over the relevant
 901 region, then small endpoint neighborhoods have small preimages under the inverse flow. Writing

$$\Lambda_T := \sup_{x \in S_0} \log |\det D_x\varphi(T, x)|$$

902 as an effective finite-time volume-expansion exponent over the relevant region, one expects an r -
 903 scale endpoint neighborhood to have mass of order $\gamma_T(r) \asymp e^{-\Lambda_T}r^n$,² up to density, anisotropy,
 904 and curvature-dependent constants. Thus an instance-dependent analysis can replace the worst-
 905 case exponent nLT by a smaller effective volume-growth exponent Λ_T , giving the heuristic sample
 906 requirement $N \gtrsim e^{\Lambda_T}r^{-n}$. In the isotropic case $\varphi(T, x) \approx e^{\lambda T}x$, we have $\Lambda_T = n\lambda T$, recover-
 907 ing the scaling $e^{n\lambda T}r^{-n}$. Hence, while the minimax theorem uses the worst-case global Lipschitz
 908 expansion nLT , persistent positive local volume expansion can still create an exponential-in-time
 909 sampling burden for Hausdorff-accurate reachable-set recovery.

910 C Experiments Details

911 This section provides additional implementation details and experimental results. We first describe
 912 the simulation platform, hardware configuration, robot-arm dynamics, parameter settings, and the
 913 estimator based on Christoffel functions used in the supplementary experiments. We then report ad-
 914 ditional results on dimension-dependent approximation of robot-arm uncertainty propagation, fol-
 915 lowed by reachable-set approximation results obtained using the Christoffel estimator.

916 C.1 Experimental Setup

917 **Simulation Platform and Hardware.** All simulations are run on a 3.2 GHz AMD Ryzen 7
 918 7735HS CPU with 16 GB RAM. The robot-arm simulations are implemented in MuJoCo 3.8.1.
 919 We use deterministic MuJoCo MJCF models with a fixed integration time step of 2×10^{-3} seconds.
 920 The approximation error is computed offline from simulated terminal samples. In the main experi-
 921 ments in Section 5, we use the convex hull of the sampled endpoints as the reachable-set estimator.

922 **Dynamics and Parameter Settings of Robotic Arms.** We consider vertical planar serial n -link
 923 robot arms with $n \in \{2, 3, 4\}$. The state is $x = [q^\top, v^\top]^\top \in \mathbb{R}^{2n}$, where $q \in \mathbb{R}^n$ denotes joint angles
 924 and $v = \dot{q} \in \mathbb{R}^n$ denotes joint velocities. Each link has length 0.5, capsule radius 0.035, density
 925 1000, joint damping 0.2, and armature 0.01. Gravity is enabled with acceleration $(0, 0, -9.81)$.
 926 MuJoCo simulates the rigid-body dynamics

$$M(q)\dot{v} + C(q, v)v + g(q) = \tau,$$

927 where $\tau \in \mathbb{R}^n$ is the vector of joint torques. For the simulation, we consider a non-adaptive inverse-
 928 dynamics tracking controller. The desired reference trajectory is

$$q_{d,i}(t) = q_{c,i} + A_i \sin(\omega t + \phi_i),$$

² $A \lesssim B$ means $A \leq cB$, $A \gtrsim B$ means $A \geq cB$, and $A \asymp B$ means both hold, for constants independent of the scale parameter.

929 with $q_c = \text{linspace}(0.25, 0.65, n)$, $A_i = 0.08$, $\omega = 0.5$, $\phi = \text{linspace}(0, \pi/3, n)$. The analytic
 930 derivatives $\dot{q}_d(t)$ and $\ddot{q}_d(t)$ are used in the controller. Let $e = q - q_d(t)$ and $\dot{e} = v - \dot{q}_d(t)$. MuJoCo
 931 inverse dynamics is used to compute

$$\tau_{\text{track}} = M(q)\ddot{q}_d(t) + C(q, v)v + g(q).$$

932 The applied torque is

$$\tau = \tau_{\text{track}} - K_p e - K_d \dot{e}.$$

933 We use weak tracking gains

$$K_p = 0.01I_n, \quad K_d = 0.005I_n,$$

934 and clip torques to $[-100, 100]$ for numerical stability. The weak gains are chosen deliberately: the
 935 benchmark is designed for uncertainty propagation and finite-sample reachable-set approximation,
 936 rather than aggressive trajectory tracking or controller design.

937 The initial uncertainty set is the box $S_0 = ([-\rho_q, \rho_q]^n \times [-\rho_v, \rho_v]^n) \circ \mathcal{B}(0, r)$ with $\rho_q = \rho_v =$
 938 0.1 , $r = 0.01$. Each trajectory is propagated to time $T = 1.0$ second, and the terminal propagated
 939 uncertainty set is approximated from sampled initial conditions.

940 **Estimator using Christoffel function.** To further validate the geometric implications of our theory,
 941 we use the empirical inverse Christoffel function as an independent verification tool [17]. Given
 942 endpoint samples $\{Y_i\}_{i=1}^N$, we construct

$$\widehat{M}_{m, \sigma_0} = \sigma_0^2 I + \frac{1}{N} \sum_{i=1}^N z_m(Y_i) z_m(Y_i)^\top,$$

943 where z_m is the vector of monomials of degree at most m , and define

$$C(y) = z_m(y)^\top \widehat{M}_{m, \sigma_0}^{-1} z_m(y).$$

944 Following the polynomial Christoffel estimator in [18, Algorithm 3], we use the sublevel set $\widehat{S}_N^{\text{Ch}} =$
 945 $\{y : C(y) \leq \eta\}$ as a data-driven certificate of support coverage. In our setting, this construction is
 946 not used as a competing reachable-set estimator, but rather as an auxiliary verification procedure: if
 947 the endpoint samples generated under our sampling scheme are sufficiently informative, the resulting
 948 Christoffel sublevel set should concentrate around the reachable set and provide an independent
 949 check of the predicted finite-sample behavior.

950 C.2 Adversarial Sampling Details

951 For completeness, we describe the adversarial sampling heuristic over S_0 used in the experi-
 952 ments to improve empirical coverage [61, 62]. Given the current endpoint cloud \mathcal{Y}^i , define
 953 $c^i := |\mathcal{Y}^i|^{-1} \sum_{Y \in \mathcal{Y}^i} Y$ and $Q^i := [|\mathcal{Y}^i|^{-1} \sum_{Y \in \mathcal{Y}^i} (Y - c^i)(Y - c^i)^\top + \lambda I]^{-1}$, where $\lambda > 0$
 954 is a regularization parameter. We then maximize the novelty objective $\mathcal{L}^i(x) := \|\varphi(T, x) - c^i\|_{Q^i}^2$,

Algorithm 1 Reachable Set Estimation via Adversarial Sampling

```

1: Input: Initial samples  $\{X_j^0\}_{j=1}^M \subset S_0$ 
2: Parameters: Step size  $\eta$ , iterations  $n_{\text{adv}}$ 
3: Output: Estimate  $\widehat{S}_N^{\text{adv}}$ 
4:  $Y_j^0 \leftarrow \varphi(T, X_j^0)$ ,  $j = 1, \dots, M$ 
5:  $\mathcal{Y}^0 \leftarrow \{Y_j^0\}_{j=1}^M$ 
6: for  $i = 0, \dots, n_{\text{adv}} - 1$  do
7:   Compute  $c^i, Q^i$  from  $\mathcal{Y}^i$ 
8:   for  $j = 1, \dots, M$  do
9:      $X_j^{i+1} \leftarrow \text{Proj}_{S_0}(X_j^i + \eta \nabla_x \mathcal{L}^i(X_j^i))$ 
10:     $Y_j^{i+1} \leftarrow \varphi(T, X_j^{i+1})$ 
11:   end for
12:    $\mathcal{Y}^{i+1} \leftarrow \mathcal{Y}^i \cup \{Y_j^{i+1}\}_{j=1}^M$ 
13: end for
14:  $N \leftarrow M(n_{\text{adv}} + 1)$ 
15: return  $\widehat{S}_N^{\text{adv}} = C(\mathcal{Y}^{n_{\text{adv}}})$ 

```

955 $x \in S_0$, which encourages new samples to generate endpoints far from the current cloud after covari-
 956 ance normalization [19]. Thus, the procedure allocates samples toward geometrically underexplored
 957 regions of the reachable set. At each iteration, the current endpoint cloud is updated along $\nabla_x \mathcal{L}^i$,
 958 projected back to S_0 , and propagated through the dynamics. The newly generated endpoints are then
 959 added to the sample cloud, and the final estimator is applied to all accumulated endpoints. For more
 960 details, please refer to [19].

961 C.3 Additional Results

962 **Fitted Dimension-Dependent Slopes.** To summarize the dimension-dependent trend in Table 2,
 963 we fit the magnitude of the empirical log–log slope as a function of the state dimension n using

$$m(n) = \frac{1}{an^b + c}, \quad \text{slope}(n) \approx -m(n).$$

964 Since the fitted curves are obtained from plots of $\log d_H$ versus $\log N$, they can also be interpreted
 965 as empirical sample-complexity exponents. Indeed, if

$$\log d_H \approx \alpha(n) - m(n) \log N,$$

966 then $d_H \approx A_n N^{-m(n)}$, where $A_n = e^{\alpha(n)}$ is the dimension-dependent prefactor in the empirical
 967 scaling law; equivalently, it is the value predicted by the fitted power law at $N = 1$. In practice, A_n
 968 should be interpreted as an intercept parameter rather than as a reliable one-sample approximation
 969 error, since the fit is obtained over a finite range of sample sizes. This prefactor captures effects
 970 not represented by the slope alone, including the geometric scale of the reachable set, flow-induced
 971 expansion, sampling density, and estimator bias. Empirically, A_n often increases with the state
 972 dimension, which further amplifies the sample requirement. This interpretation is consistent with
 973 classical support- and level-set estimation results, where Hausdorff-type recovery rates and covering
 974 complexities depend explicitly on the ambient or intrinsic dimension [63]. Thus achieving accuracy
 975 $d_H \leq r$ requires, up to dimension-dependent constants, $N \gtrsim \left(\frac{A_n}{r}\right)^{1/m(n)} = \left(\frac{A_n}{r}\right)^{an^b+c}$. There-
 976 fore, the fitted denominator $an^b + c$ represents the empirical exponent with which the sample budget
 977 must grow as the target accuracy r decreases.

978 The corresponding parameters and fitted curves are shown in Table 3 and Figure 7, respectively.

979 Substituting the fitted parameters in Table 3 into the relation $N \gtrsim \left(\frac{A_n}{r}\right)^{an^b+c}$ shows that the sample
 980 budget grows as

$$N_{\text{unif}} \gtrsim \left(\frac{A_n}{r}\right)^{0.5049 n^{1.0709} + 1.3353}, \quad N_{\text{adv}} \gtrsim \left(\frac{A_n}{r}\right)^{0.9488 n^{0.7203} + 0.2535}.$$

Table 3: Fitted parameters for the empirical slope model $m(n) = 1/(an^b + c)$, where n is the state dimension and slope $(n) \approx -m(n)$. The implied sample-complexity exponent is $1/m(n) = an^b + c$.

Sampling method	a	b	c
Uniform sampling	0.5049	1.0709	1.3353
Adversarial sampling	0.9488	0.7203	0.2535

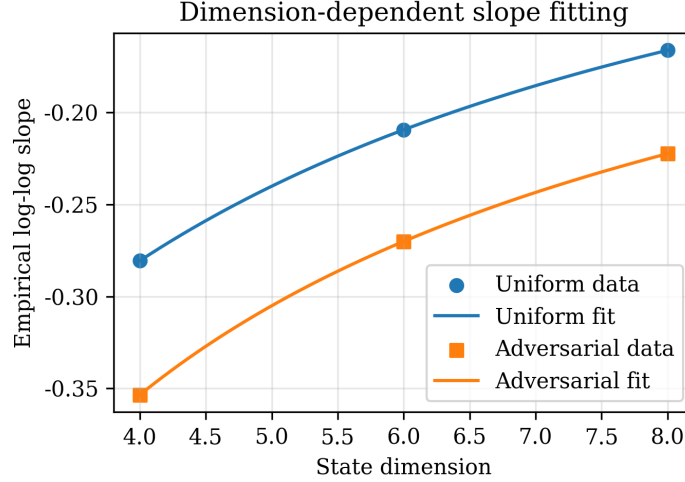


Figure 7: Fitted dimension-dependent empirical slopes for robot-arm uncertainty propagation. The fitted model captures the flattening of the log-log convergence slopes as the state dimension increases.

981 Equivalently, taking logarithms on both sides yields

$$\log N_{\text{unif}} \gtrsim (0.5049 n^{1.0709} + 1.3353) \log \frac{A_n}{r}, \quad \log N_{\text{adv}} \gtrsim (0.9488 n^{0.7203} + 0.2535) \log \frac{A_n}{r}.$$

982 Thus the empirical sample requirement does not grow only logarithmically with dimension. Both
 983 uniform and adversarial sampling exhibit an exponential-type dependence on the state dimension
 984 through the exponent multiplying $\log(A_n/r)$. Adversarial sampling reduces the fitted exponent,
 985 but the exponent still increases with n , indicating that targeted sampling improves finite-sample
 986 efficiency without eliminating the intrinsic curse of dimensionality.

987 **Sample-Budget-Dependent Improvement of Adversarial Sampling.** Section 5 reports the main
 988 dimension-dependent robot-arm results. Here we provide an additional comparison showing how
 989 the benefit of adversarial sampling depends on the sample budget. Table 4 reports the absolute and
 990 relative improvement of adversarial sampling over uniform sampling for the 2-, 3-, and 4-link robot
 991 arms.

992 The advantage of adversarial sampling is not uniform across all sample sizes. When N is very small,
 993 adversarial updates may over-concentrate samples along a few extreme directions and reduce global
 994 coverage. As the sample budget grows, the sampler has enough coverage to exploit informative
 995 directions, and the relative improvement becomes positive and more stable. At $N = 3000$, the
 996 relative improvements are approximately 39.09%, 35.56%, and 30.58% for the 2-, 3-, and 4-link
 997 arms, respectively.

998 **Time dependence under closed-loop control.** Figure 8 shows that, under the closed-loop robot-
 999 arm dynamics, the Hausdorff error grows only slowly with the simulation horizon T . This contrasts
 1000 with the non-Lipschitz and Linear system examples in Section 5.1, where flow expansion strongly
 1001 amplifies sampling errors. In the robot-arm setting, the inverse-dynamics tracking controller reduce

Table 4: Absolute (Abs.) and relative (Rel.) improvement of adversarial sampling over uniform sampling for robot-arm uncertainty propagation.

N	$n = 2$		$n = 3$		$n = 4$	
	Abs.	Rel. (%)	Abs.	Rel. (%)	Abs.	Rel. (%)
1	-0.0117	-3.35	0.0000	0.01	-0.0098	-2.15
3	0.0236	8.27	-0.0157	-4.88	-0.0126	-3.47
10	0.0356	17.79	0.0122	4.93	0.0011	0.37
30	0.0365	25.86	0.0374	18.01	0.0250	10.26
100	0.0397	36.72	0.0316	20.70	0.0386	18.84
300	0.0317	40.10	0.0307	25.35	0.0462	25.90
1000	0.0226	41.17	0.0290	30.12	0.0371	26.68
3000	0.0144	39.09	0.0280	35.56	0.0355	30.58

1002 the effective expansion of the uncertainty set, so the observed time dependence is much milder than
 1003 the worst-case exponential growth predicted by the minimax bound.

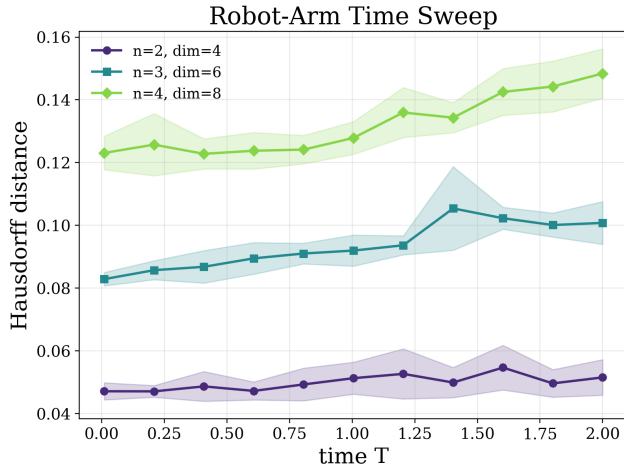


Figure 8: Hausdorff error versus time under uniform sampling

1004 **Effect of Adversarial Sampling Intensity Across Estimators.** We further study how the adver-
 1005 sarial intensity of the sampling procedure affects reachable-set approximation under different down-
 1006 stream estimators. We repeat the time-sweep experiment for the autonomous system $\dot{y} = 0, \dot{x} = x^2$
 1007 using the same three initial sets as in Section 5.1: a disk, a triangle, and an opened triangle with
 1008 equal area. We compare two estimators constructed from the same endpoint samples: the convex
 1009 hull estimator and a Christoffel-type estimator.

1010 The adversarial intensity is controlled by the number of adversarial updates $n_{\text{adv}} \in \{0, 1, 2, 3, 4\}$,
 1011 where $n_{\text{adv}} = 0$ corresponds to purely uniform sampling. Larger n_{adv} allocates a larger fraction
 1012 of the fixed sample budget to points obtained after adversarial updates, and therefore places more
 1013 emphasis on expanding or boundary-like directions in the endpoint cloud.

1014 Figures 9 and 10 show the Hausdorff error over time for sample budgets $N = 10, 100, 1000$. The re-
 1015 sults show that the effect of adversarial sampling is both budget-dependent and estimator-dependent.
 1016 For very small sample budgets, increasing n_{adv} does not consistently improve performance, since
 1017 overly concentrated adversarial samples may reduce global coverage. For larger budgets, adversarial
 1018 updates become more beneficial: the endpoint cloud already has enough global coverage, and addi-
 1019 tional adversarial samples help resolve boundary and extreme regions. This trend is visible for both
 1020 convex-hull and Christoffel estimators, although the best choice of n_{adv} varies across estimators,
 1021 sample budgets, and initial geometries. Overall, these results suggest a coverage–boundary tradeoff:

1022 stronger adversarial sampling can improve finite-sample performance, but only when the sample
 1023 budget is large enough to avoid sacrificing global coverage.

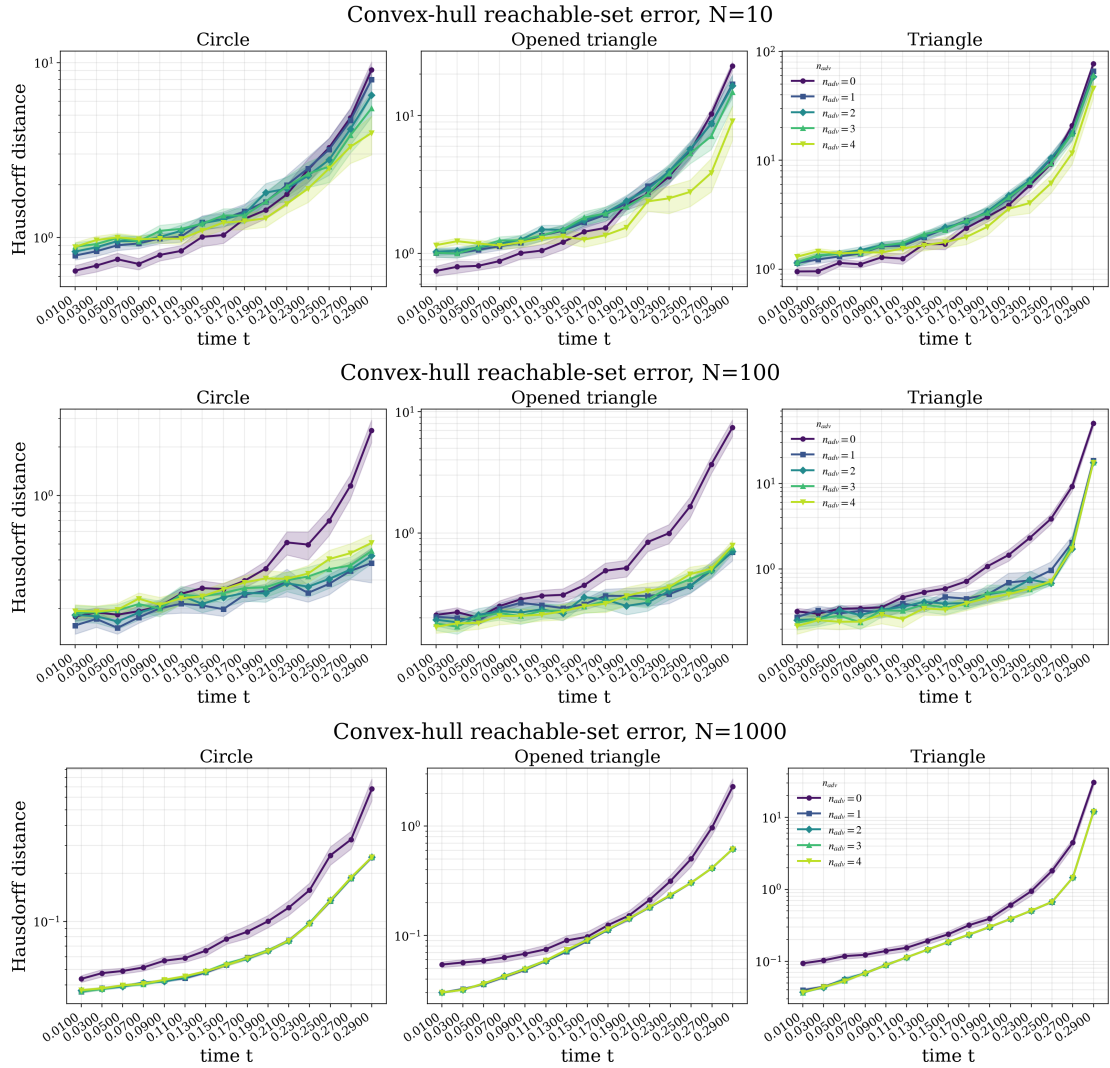


Figure 9: Hausdorff error versus time for Convexhull reachable-set approximation of the autonomous system $\dot{y} = 0, \dot{x} = x^2$ under different uniform-adversarial mixture sampling ratios. The three subfigures correspond to sample budgets $N = 10, 100, 1000$, and each subfigure compares the disk, triangle, and opened-triangle initial sets.

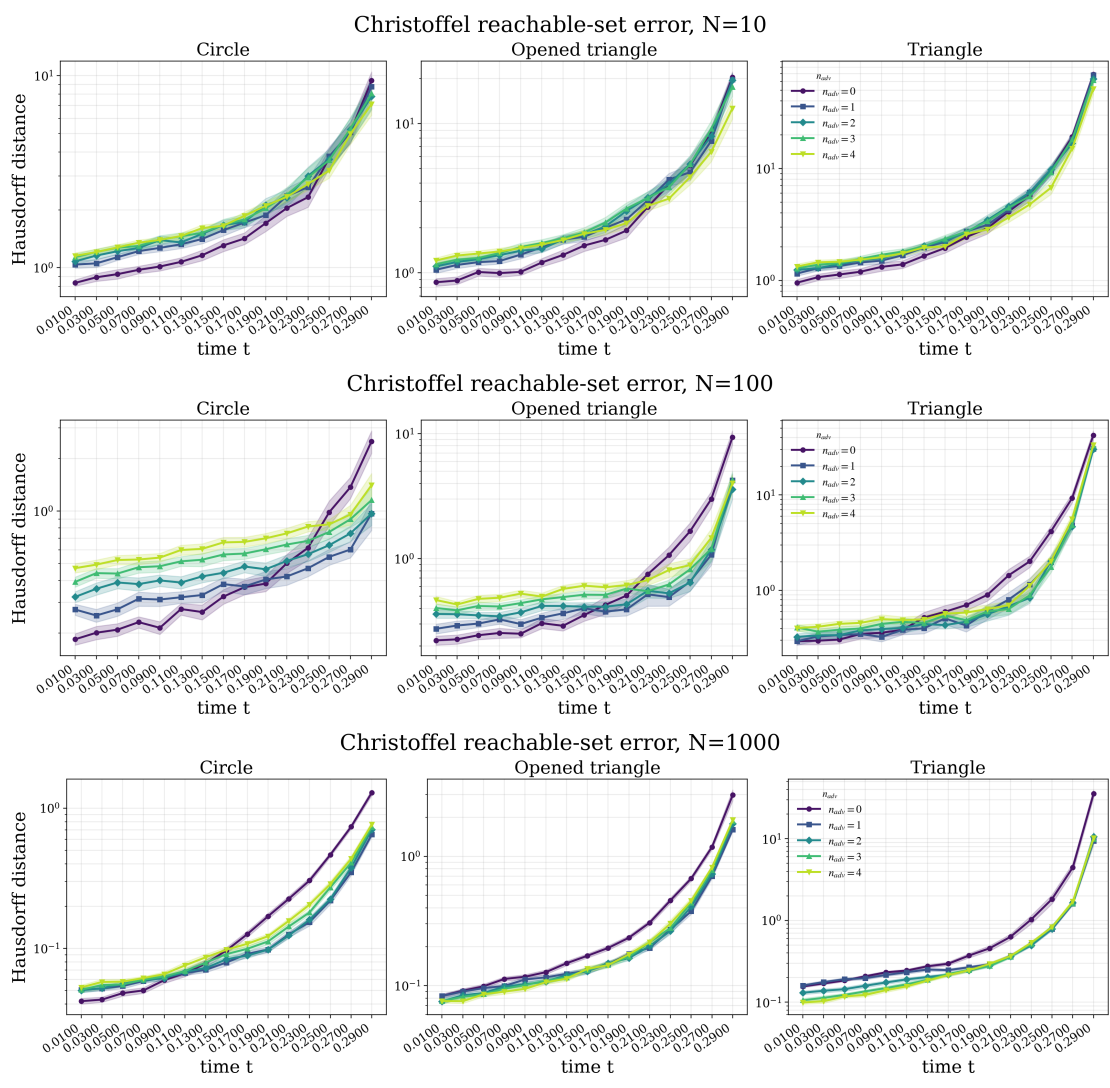


Figure 10: Hausdorff error versus time for Christoffel reachable-set approximation of the autonomous system $\dot{y} = 0, \dot{x} = x^2$ under different uniform-adversarial mixture sampling ratios. The three subfigures correspond to sample budgets $N = 10, 100, 1000$, and each subfigure compares the disk, triangle, and opened-triangle initial sets.